### Salami tactics

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### The basic idea

- This design technique can be applied for search problem
- We are given a set S. We want to find a specific element of it, or verify that S doesn't contain that element.
- Instead of locating the goal element, we "fast" exclude as "many" elements from the search space as possible (we must guarantee that the thrown away elements don't contain the goal element).
- The best way to understand this scheme is to see examples.

# The scheme by picture



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#### Binary search

Given the ordered sequence of *n* real numbers:  $x_1 < x_2 < ... < x_n$ . Also given a number *g*. Decide whether *g* is among our  $x_i$ 's. If yes, then provide the index *i* for that  $x_i = g$ .

We assume that  $n = 2^k - 1$ . Specially in the ordered sequence there is a median/middle element:  $x_{2^{k-1}}$ .

## The algorithm

#### Binary search algorithm

Given the ordered sequence of *n* real numbers:  $x_1 < x_2 < \ldots < x_n$ , where  $n = 2^k - 1$ , and a number, *g*: the "goal".

(0) If n = k = 1, after comparing our only number and g we can answer the problem. Otherwise

The above algorithm uses k comparisons if g is not among our numbers. If g is one of the  $x_i$ -s it might stop earlier.



# A popular form of binary search

- We have s, a secret number from the set  $\{1, 2, 3, \dots, 99, 100\}$ .
- By asking Boolean questions about s (yes/no answer) figure out the value of s.
- How many questions are needed to achieve this goal?

### Break



### The basic problem

#### Definition: The median problem

Given *n* numbers:  $x_1, x_2, \ldots, x_n$ . Find the index *m*:  $x_m$  is the median of our numbers.

# A generalization

### Definition: Generalized median problem

Given *n* numbers:  $x_1, x_2, ..., x_n$ , and a rank number  $r \in \{1, 2, ..., n-1, n\}$ . Classify the set of indices into three classes,  $S \cup \{m\} \cup B = \{1, 2, ..., n-1, n\}$ , the following way: (1) for each  $s \in S$  we have  $x_s \leq x_m$ , (2) for each  $b \in B$  we have  $x_m \leq x_b$ , (3) |S| = r - 1. // Hence |B| = n - r.

Note that the median problem is the case when  $r = \lfloor (n+1)/2 \rfloor$ , furthermore we are not interested in the sets S and B.

Again we have a simplifying assumption: *n* has the form  $10k + 5 = (2k + 1) \cdot 5$ 

# The algorithm

### Generalized median finding algorithm

Given *n* numbers, and a rank *r*. // Think about the *n* input numbers as they are arranged in a matrix of size  $5 \times (2k + 1)$ .

(Column medians) For each columns (number quintets) determine the median and their S, B sets. ( $m_i$  is the median of the *i*th column, furthermore  $S_i$ ,  $B_i$  are the corresponding sets.)

(Median of medians) For the n/5 = 2k + 1 medians find the median:  $\mu$ , and the corresponding sets:  $\tilde{S}$ ,  $\tilde{B}$ . // Note the following set contains only numbers that are not bigger than  $\mu$ :

$$\widehat{S} = \bigcup \{\{m_i\} \cup S_i : m_i \in \widetilde{S}\} \cup S_i : m_i = \mu.$$

// The following set contains only numbers that are not smaller than  $\mu$ :

$$\widehat{B} = \bigcup \{\{m_i\} \cup B_i : m_i \in \widetilde{B}\} \cup \{B_i : m_i = \mu\}.$$

# The algorithm (cont'd)

### Generalized median finding algorithm (cont'd)

(Finding the rank of  $\mu$ ) Compare  $\mu$  to all other elements. Determine the corresponding  $S_{\mu}$ ,  $B_{\mu}$  sets.

 $ho = |S_{\mu}| + 1.$ // We have  $\widehat{S} \subset S_{\mu}$ , and  $\widehat{B} \subset B_{\mu}$ .

(The cut of the salami) If  $\rho = r$ , then we are extremely lucky. If  $\rho < r$ , then throw away the numbers of  $\widehat{S}$ .  $r \leftarrow r - |\widehat{S}|$ . If  $\rho > r$ , then throw away the numbers of  $\widehat{B}$ .

(Iteration) If the cardinality of the leftover numbers is smaller or equal to five then solve the problem by brute force.

Otherwise do the same for the leftover numbers, and r.

# A claim

### Claim

Let  $n = (2k + 1) \cdot 5$  as in the algorithm. Then

$$|\widehat{S}| = |\widehat{B}| = 3k + 2 = 3n/10 + 1/2.$$

### Corollary

Let  $n = (2k + 1) \cdot 5$  as in the algorithm. Then in the first iteration we reduced the input size to

$$n - 3k + 2 = n - (3n/10 + 1/2) = 7n/10 - 1/2.$$

### The analysis of the algorithm: Introduction

Let M(n) be the maximal number of comparisons needed for our algorithm to solve the problem on n numbers.

 $M(x) = M(\lceil x \rceil).$ 

Observation

M(x) is a monotone increasing function.

# The analysis of the algorithm

#### Theorem

### $M(n) \le M((n+9)/5) + M(7n/10) + O(n)$

The third term counts the comparisons needed for (Column medians) and for (Finding the rank of  $\mu$ ).

#### Theorem

$$M(n) \leq M(0.21 \cdot n) + M(0.7 \cdot n) + O(n),$$

assuming n > 200

### Corollary

$$M(n)=O(n),$$

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# Thank you for your attention!