

Geometric constructibility of cyclic polygons and a limit theorem

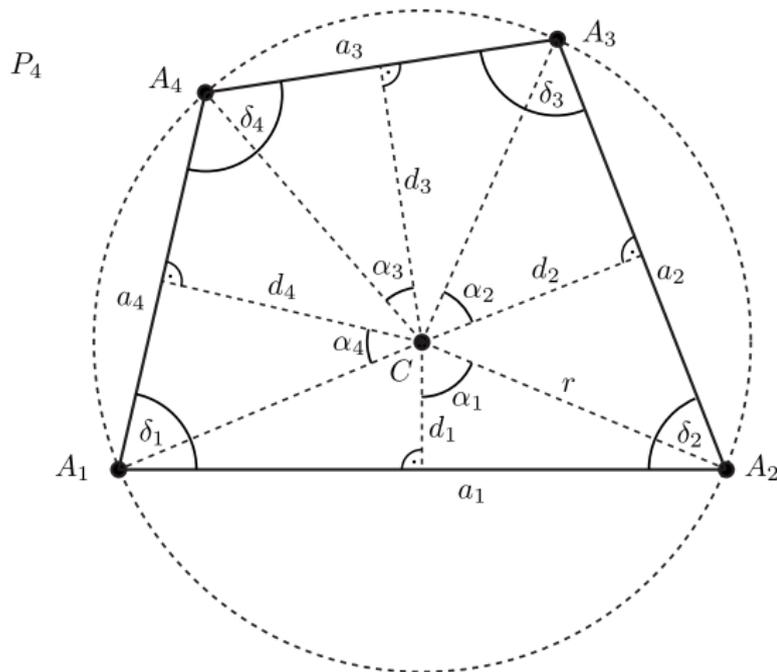
Gábor Czédli and **Ádám Kunos**
(SSAOS, Srní, August 29–September 4, 2015)

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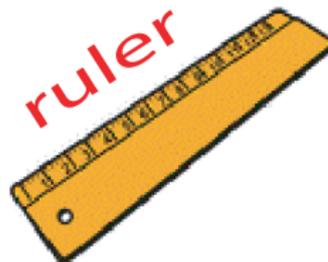
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September 3, 2015



Cyclic n -gon $\stackrel{\text{def}}{\iff}$
 all of its ver-
 tices are on the
 same circle (AKA
 inscribed polygon)

Constructibility := **geometric constructibility** with straight-edge and compass (“Straight-edge” is more precise than “ruler”.)



Schreiber, P.: On the existence and constructibility of inscribed polygons. *Beiträge zur Algebra und Geometrie* **34**, 195–199 (1993)

Theorem (Schreiber): The cyclic pentagon cannot be constructed from its side lengths.

He also claims: for every $n > 5$, the cyclic n -gon cannot be constructed from its side lengths.

13 line long argument **wrong**

In essence: Suppose P_n is constructible for some $n > 5$. Let $a_n \rightarrow 0$, then $P_n \rightarrow P_{n-1}$, so P_{n-1} is also constructible. Etc., $P_6 \rightarrow P_5$, so the cyclic pentagon is constructible, a contradiction.

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Theorem (Czédli–Kunos, 2015)

- For $n \geq 5$, the cyclic n -gon is not constructible from its side lengths in general.
- Furthermore, for each $n \geq 5$, there exist positive integers a_1, \dots, a_n such that the cyclic n -gon with these side lengths exists, it is not constructible from a_1, \dots, a_n , and

$$|\{a_1, \dots, a_n\}| = \begin{cases} 2, & \text{if } n \neq 6, \\ 3, & \text{if } n = 6. \end{cases}$$

For $n = 5$: Schreiber, 1993. For $n \in \{6, \dots, 100\}$: Czédli and Ágnes Szendrei, 1997.

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$K(\sqrt{c}) := \{a + b\sqrt{c} : a, b \in F\}$ simple quadratic field extension

$K = K_0 \subset K_1 \subset \dots \subset K_m$ quadratic tower

K^\square is the smallest subfield of \mathbb{R} such that includes K and is closed with respect to real square roots. real quadratic closure

Known: b is constructible from a_1, \dots, a_n iff b is a real quadratic number over $\mathbb{Q}(a_1, \dots, a_n)$, that is, iff $b \in \mathbb{Q}(a_1, \dots, a_n)^\square$.

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quadratic
tower of
length k
over $F(x)$

Its elements, like f , are just expressions (=strings).

Always, $F \subseteq \mathbb{R}$. $*f$: the real function determined by $f \in F_k\langle x \rangle$ with the largest possible domain (induction!), $\text{Dom}(*f) \subseteq \mathbb{R}$.

Note that $f \mapsto *f$ is not injective! For example,

$\text{Dom}(*(\sqrt{-1-x^2})) = \emptyset = \text{Dom}(*(\sqrt{-1-x^4}))$ and

$*(\sqrt{-1-x^2}) = *(\sqrt{-1-x^4})$, but $\sqrt{-1-x^2} \neq \sqrt{-1-x^4}$.

Corollary (Corollary 6.4 in the paper)

If $|\text{Dom}(*f_1) \cap \text{Dom}(*f_2)| = \infty$ and $f_1 \neq f_2$, then $*f_1 \neq *f_2$.

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Theorem (on “quadratic” Puiseux series (the **heart** of the paper))

With $F \subseteq \mathbb{R}$, closed w.r. to real quadratic roots, and f as before (and above), assume $\exists(0, \varepsilon) \subseteq \text{Dom}(*f)$. Then there exist an integer $t \in \mathbb{Z}$ and elements $b_t, b_{t+1}, b_{t+2}, \dots \in F$ such that

$$*f(x) = \sum_{j=t}^{\infty} b_j \cdot x^{j \cdot 2^{-k}} \quad (1)$$

holds in some strict right neighborhood of 0. Furthermore, if $f \neq 0$, then $b_t \neq 0$ and $t, b_t, b_{t+1}, b_{t+2}, \dots$ are uniquely determined.

We obtain (1) from a Laurant series with finitely many powers of negative exponents, by substituting the complex variable z with $\sqrt[k]{x}$. Using an arbitrary $\ell \in \mathbb{N}$ rather than 2^k , we get the concept of a *Puiseux series*, which goes back to Isaac Newton.

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Proof of the “quadratic version of the Newton–Puiseux theorem”.

Induction on k (looking for the b_j as unknowns) gives that $\exists b_j \in F$. (The problem reduces to square rooting, because $+$ and \cdot create no problem. At present, only in the ring of formal Puiseux series.)

Substituting $x := y^{2^k}$ and using the uniqueness of Laurent expansions \Rightarrow if we have convergence, then the b_j are uniquely determined.

For convergence: some arguments on holomorphic functions and using the binomial series, namely, for $z \in (-1, 1)$, we have

$$\sum_{j=0}^{\infty} \binom{1/2}{j} x^j = 1 + \frac{1}{2}x + \frac{\frac{1}{2}(\frac{1}{2}-1)}{2!}x^2 + \frac{\frac{1}{2}(\frac{1}{2}-1)(\frac{1}{2}-2)}{3!}x^3 + \dots$$

No time for the complicated details (why \mathbb{C} for a while, etc.). \square

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Armed with our quadratic Puiseux series, **yes**. (To be precise, **yes** in the 33 page long paper **almost** here where some issues are not discussed; they will be written in **red**.)

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Proving part •: for $n \geq 5$, the cyclic n -gon is not constructible.

Take the smallest $n > 5$ such that the cyclic n -gon is constructible. Pick $a_1, \dots, a_{n-1} \in \mathbb{R}$ such that the cyclic $(n-1)$ -gon is non-constructible. Let $x = a_n \rightarrow 0$ through irrationals numbers. $F := \mathbb{Q}(a_1, \dots, a_{n-1})^\square$. Since the cyclic n -gon is constructible “over $F(x)$ ” and so is its radius, this radius is of the form ${}^*f(x)$ for some $f \in F_k\langle x \rangle$. (Pigeon hole principle $\Rightarrow k$ and f do not depend on x .) By our “quadratic Newton–Puiseux”, ${}^*f = \sum_{j=t}^{\infty} b_j \cdot x^{j \cdot 2^{-k}}$; with the $b_j \in F$. Since ${}^*f(x) = {}^*f(a_n)$ is the radius of the n -gon that tends to the nonzero radius of the $(n-1)$ -gon, $\lim_{x \rightarrow 0+0} {}^*f(x)$ exists (and it is positive). Hence, $t = 0$, and this limit is $b_t = b_0 \in F$. But this limit is the radius of the $(n-1)$ -gon. Thus, this radius and the cyclic $(n-1)$ -gon is constructible, a contradiction. \square

Thank you for your attention.

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