

An association rule motivated by

Grätzer and Schmidt

by **Gábor Czédli**

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but

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the best result I will recall is due to

Kalle Kaarli

2006. június 8.

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- Schmidt: my PhD advisor (“Candidate Degree” 1984)
- [Schmidt, E. Tamás](#) Remarks on dependence relations in relational database models. (Hungarian. English summary) [Alkalmaz. Mat. Lapok](#) 8 (1982), no. 1-2, 177–182. [68H05 \(68B15\)](#)

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K. Kaarli: Finite uniform lattices are congruence permutable: *Acta Sci. Math. (Szeged)*, to appear.

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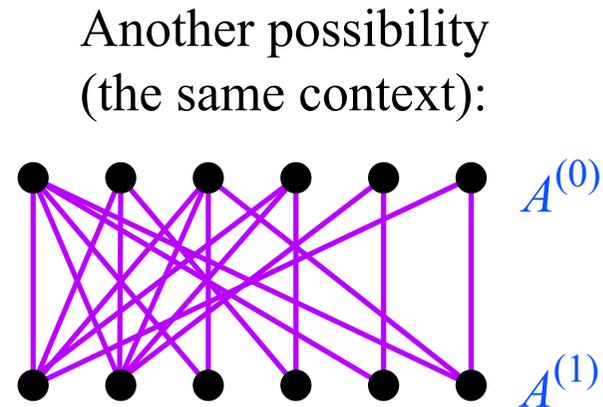
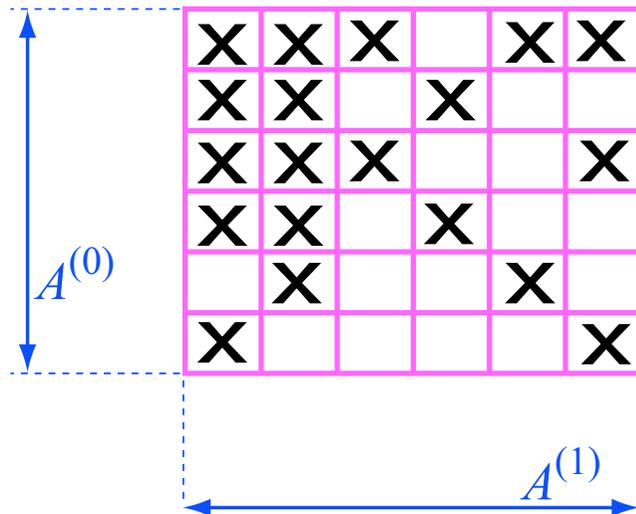
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 $(A^{(0)}, A^{(1)}, \rho)$ where $A^{(0)} = \{\text{objects}\}$, $A^{(1)} = \{\text{attributes}\}$,
 $\rho \subseteq A^{(0)} \times A^{(1)}$.



Permanent notations for $(A^{(0)}, A^{(1)}, \rho)$ and $X \in P(A^{(i)})$:

$$\begin{aligned}\rho_0 &= \rho & \rho_1 &:= \rho^{-1}, \\ X\rho_i &:= \{y \in A^{(1-i)} : \forall x \in X, (x, y) \in \rho_i\}, \\ \mathcal{C}_0^{(i)} &: P(A^{(i)}) \rightarrow P(A^{(i)}), \quad \mathcal{C}_0^{(i)}(X) := (X\rho_i)\rho_{1-i},\end{aligned}$$

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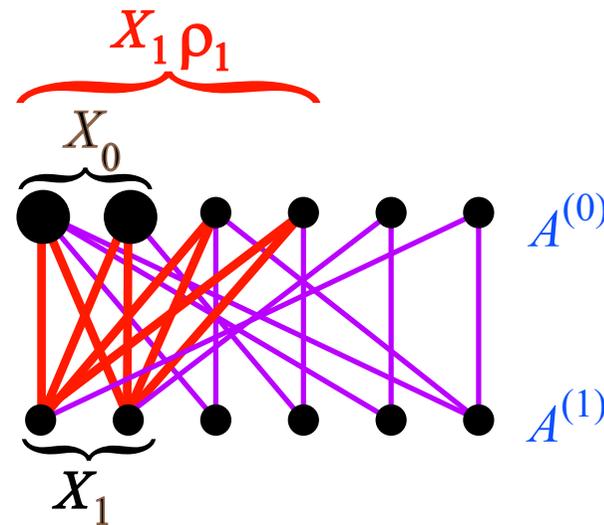
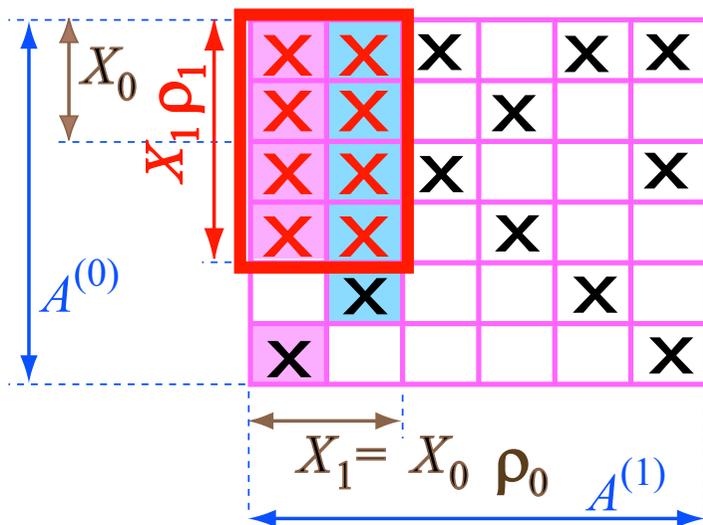
To express that $b \in \mathcal{C}_0^{(i)}(X)$, **data miners** say that $X \rightarrow b$ is a **strong association rule**. *Association rule* is the fuzzy version.

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Pictorially, $\mathcal{C}_0^{(i)}(X_0)$, the G–W closure of $X_0 \in P_0(A^{(0)})$, is defined by maximal full rectangles, in other terminology (cf. Wille), by **concepts**:



$$c_0^{(i)}(X_0) = \bigcap_{X_0 \subseteq S, S \text{ is a side of a max. full rectangle}} S.$$

Instead of maximal full rectangles we need something else.

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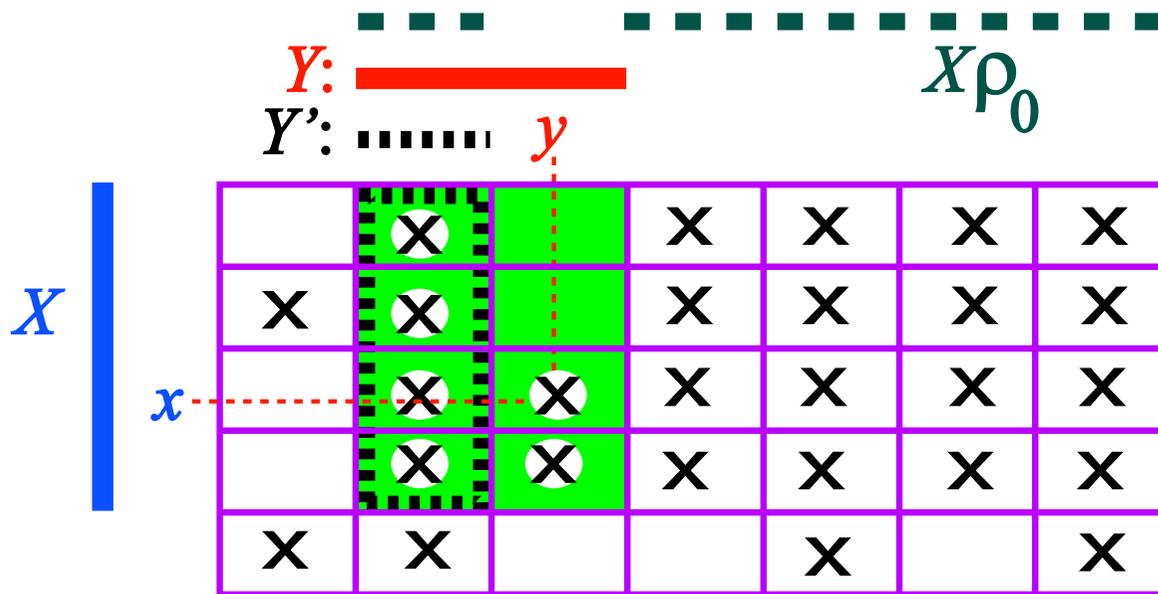
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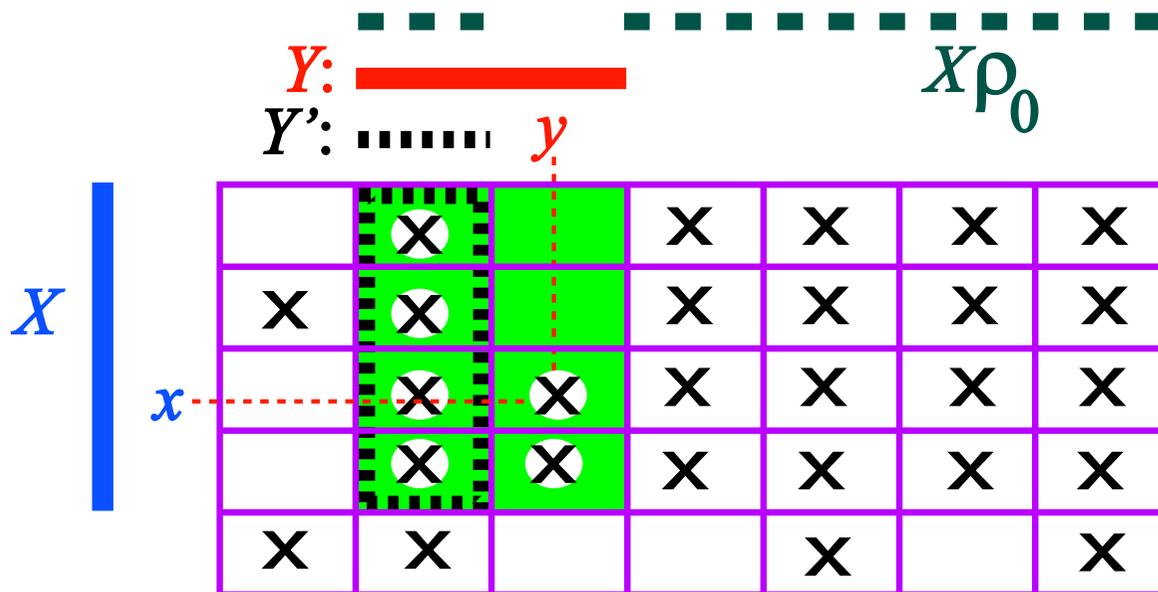
and let

$$\psi_{\min}(X) := \{Y : Y \text{ is minimal in } (\psi(X), \subseteq) \}.$$

Pictorially:



Here $Y \in \psi(X)$, $Y' \in \psi_{\min}(X)$ but



Here $Y \in \psi(X)$, $Y' \in \psi_{\min}(X)$ but $X\rho_0 \notin \psi(X)$. So the terminology „ ρ -image” should not be confused with $X\rho_0$!

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$$\mathcal{C}_{n+1}^{(i)}(X) := \mathcal{C}_n^{(i)}(X) \cap \bigcap_{Y \in \psi(X)} \bigcup_{y \in \mathcal{C}_n^{(1-i)}(Y)} \{y\} \rho_{1-i}.$$

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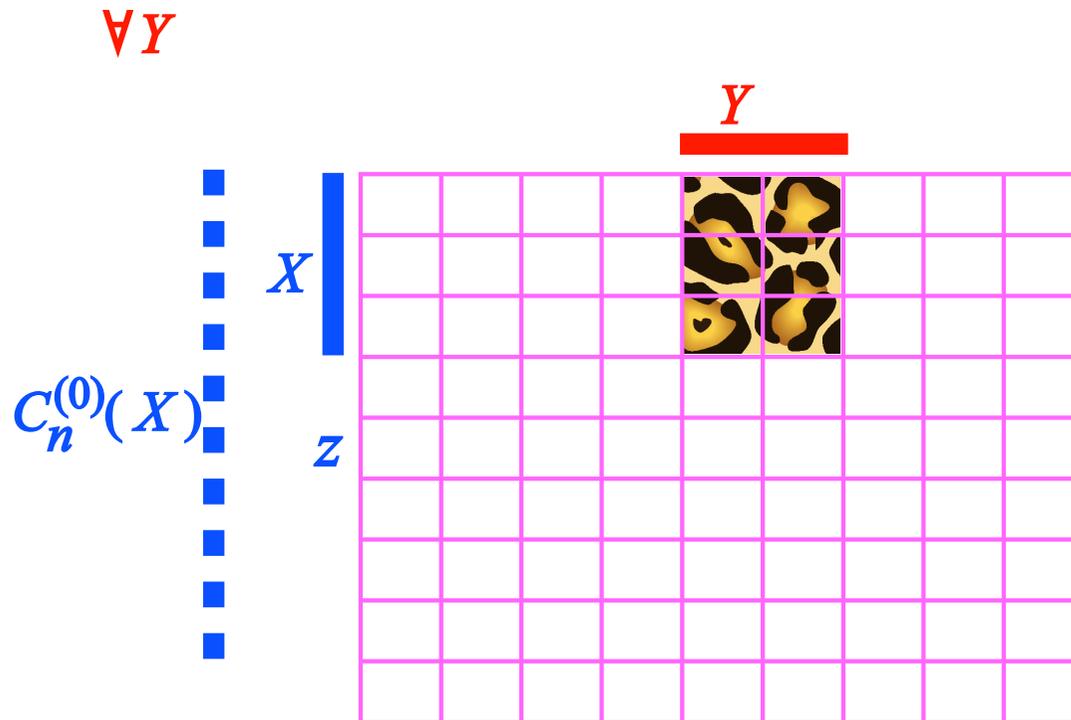
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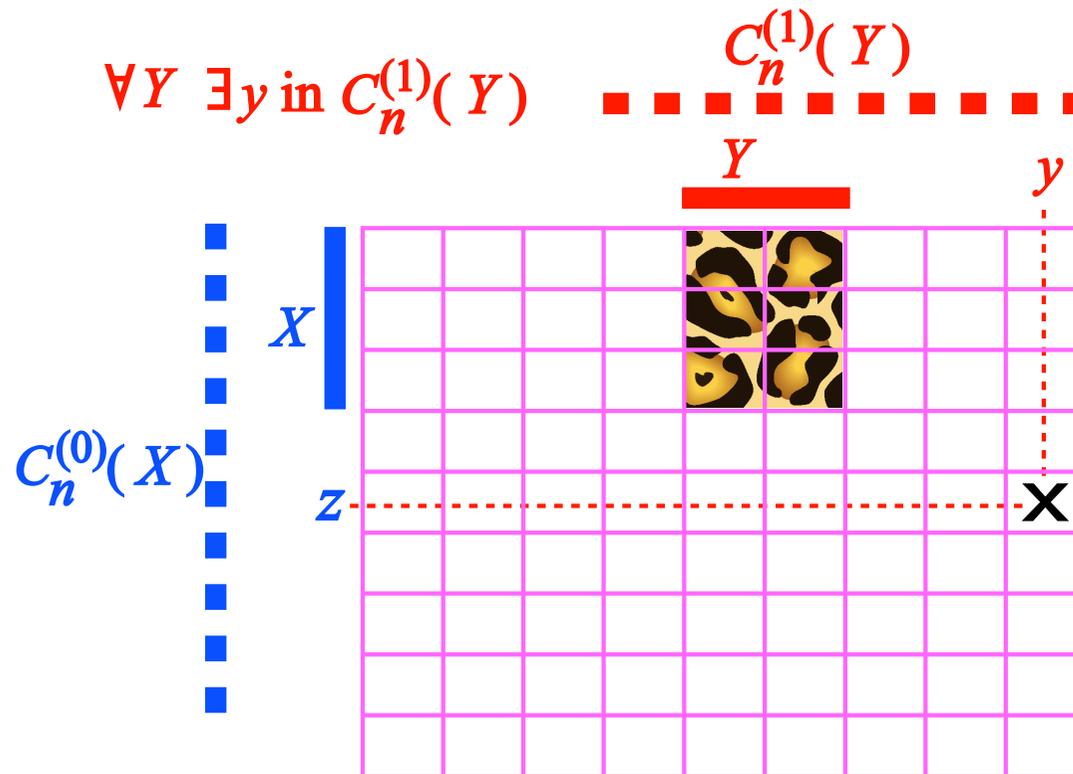
In the finite case $\psi(X)$ can be replaced by $\psi_{\min}(X)$.

For $i = 0$ we explain this formula pictorially for the finite case.
(Not all the crosses will be indicated.)

Let $z \in \mathcal{C}_n^{(0)}(X)$. Then $z \in \mathcal{C}_{n+1}^{(0)}(X)$ iff



for each $Y \in \psi_{\min}(X)$



there is a $y \in C_n^{(1)}(Y)$ with the cross indicated (i.e., $(z, y) \in \rho$).

This

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$$\mathcal{C} = (\mathcal{C}^{(0)}, \mathcal{C}^{(1)}) := \left(\bigcap_{n=0}^{\infty} \mathcal{C}_n^{(0)}, \bigcap_{n=0}^{\infty} \mathcal{C}_n^{(1)} \right),$$

which means that, for all $X \in P(A^{(i)})$,

$$\mathcal{C}^{(i)}(X) = \bigcap_{n=0}^{\infty} \mathcal{C}_n^{(i)}(X).$$

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(understood componentwise).

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Sketch of the proof. Pictorially. Suppose \mathcal{C}_n is already a pair of closure operators and we already know that $\mathcal{C}_{n+1}^{(0)}$ and $\mathcal{C}_{n+1}^{(1)}$ are **monotone** and **extensive** mappings.

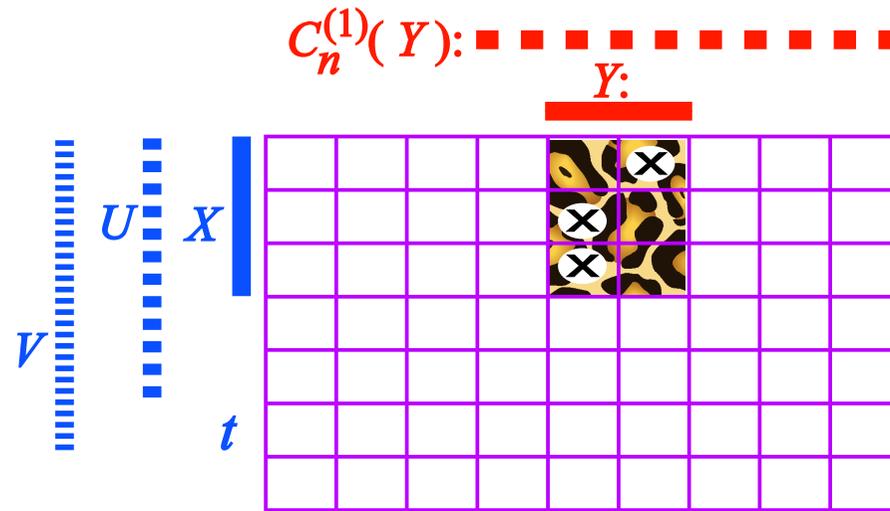
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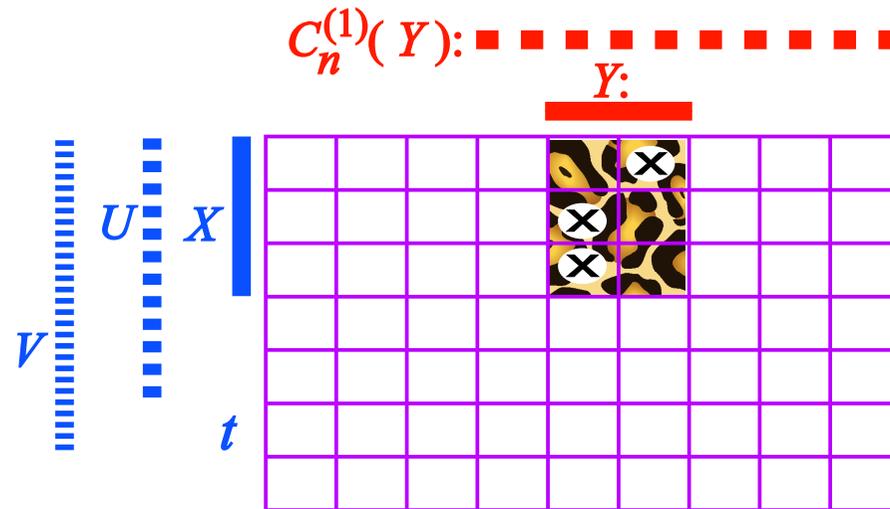
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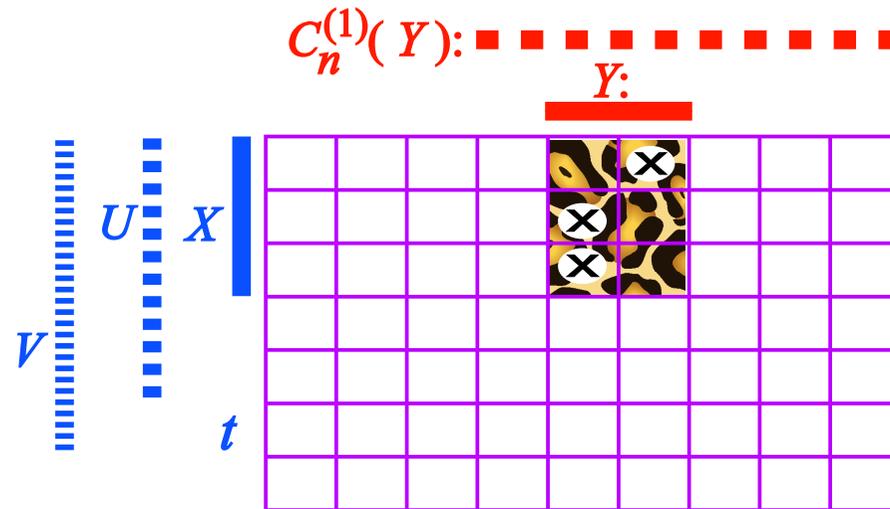
We want: $\mathcal{C}_{n+1}^{(i)}$ is **idempotent**. Let $i = 0$.



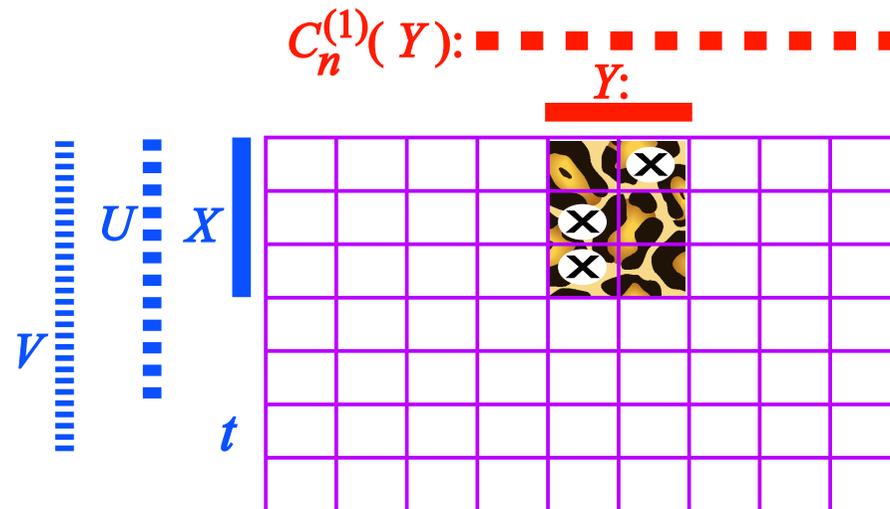
Let $X \in P(A^{(0)})$, $U = c_{n+1}^{(0)}(X)$ and $V = c_{n+1}^{(0)}(U)$.



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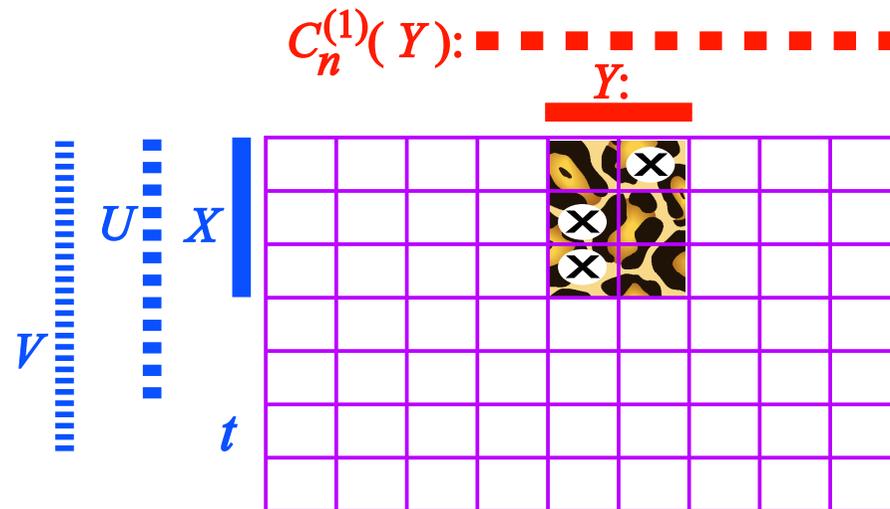


Let $X \in P(A^{(0)})$, $U = c_{n+1}^{(0)}(X)$ and $V = c_{n+1}^{(0)}(U)$. We want to show $V \subseteq U$. Suppose this is false. Then $\exists t \in V \setminus U$.

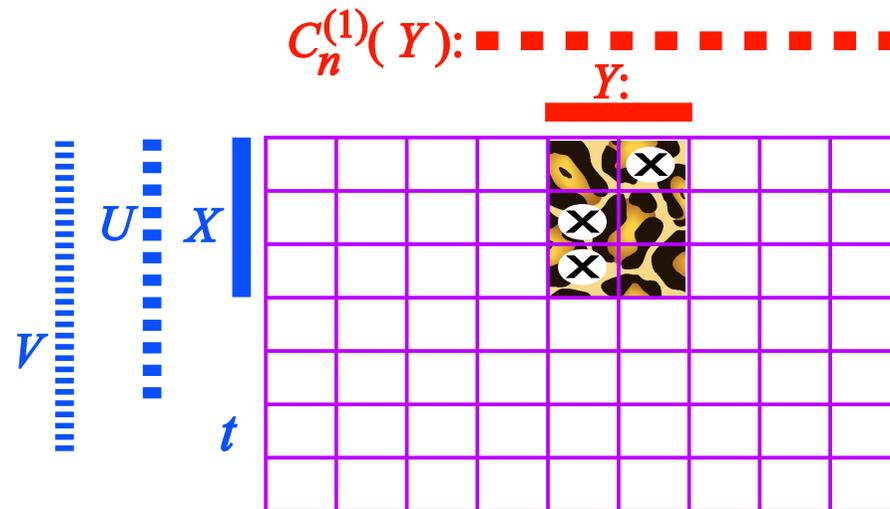


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$$U = C_{n+1}^{(0)}(X) := C_n^{(0)}(X) \cap \bigcap_{Y \in \psi(X)} \bigcup_{y \in C_n^{(1)}(Y)} \{y\} \rho_1$$

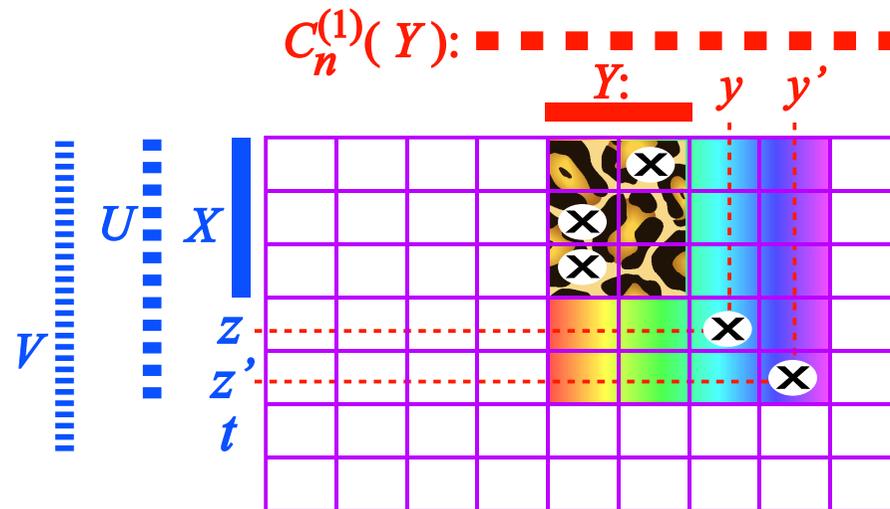


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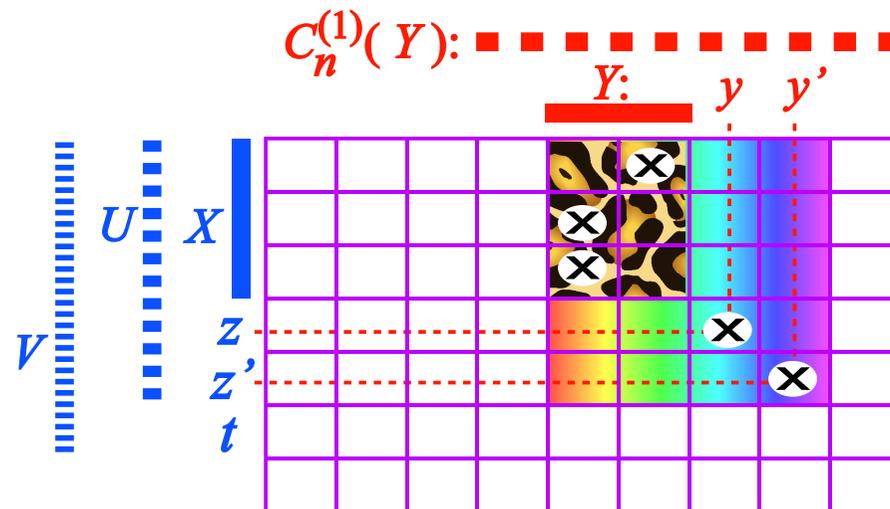
For any $Y \in \psi(X)$ we need a cross in $\{t\} \times \mathcal{C}_n^{(1)}(Y)$.

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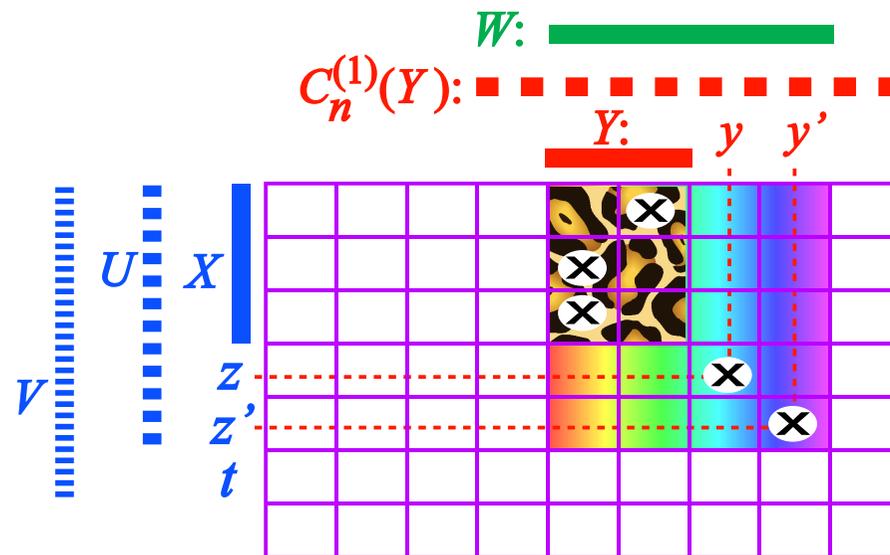
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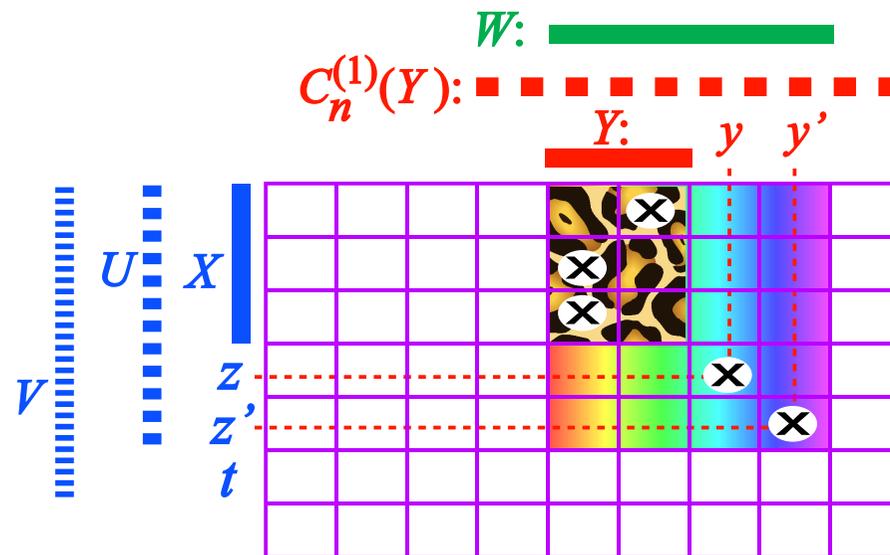
$\forall z, z', \dots \in U \setminus X \exists y, y', \dots \in \mathcal{C}_n^{(1)}(Y)$ by the def. of $U = \mathcal{C}_{n+1}^{(0)}(X)$.

$$c_{n+1}^{(0)}(X) := c_n^{(0)}(X) \cap \bigcap_{Y \in \psi(X)} \bigcup_{y \in c_n^{(1)}(Y)} \{y\} \rho_1$$



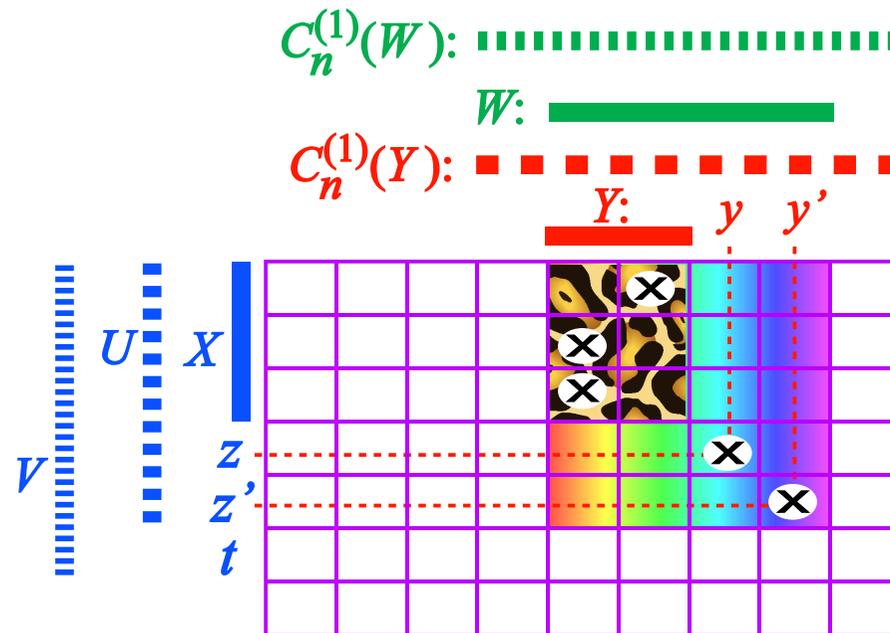
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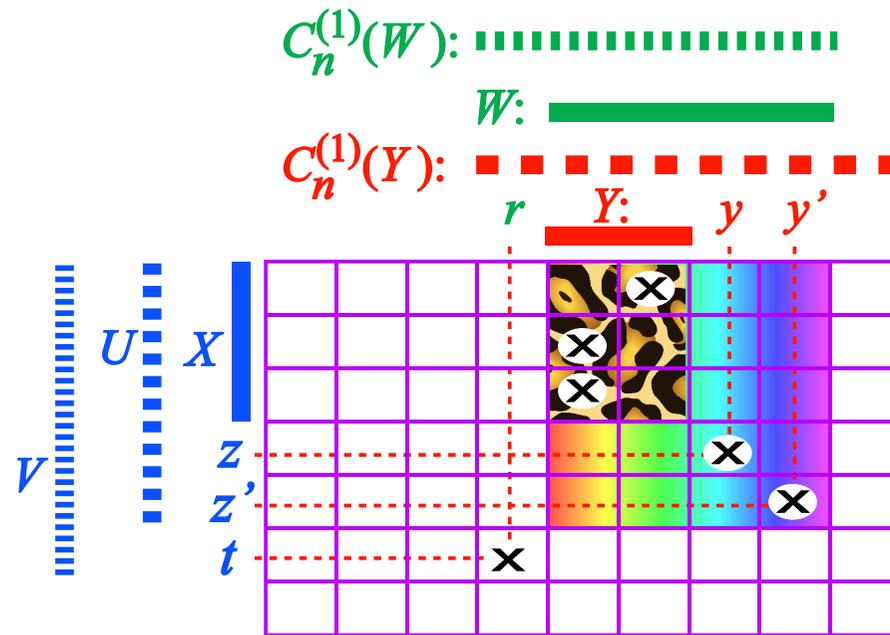
$$W := Y \cup \{y, y', \dots\}$$

$$V = c_{n+1}^{(0)}(U) := c_n^{(0)}(U) \cap \bigcap_{W \in \psi(U)} \bigcup_{r \in c_n^{(1)}(W)} \{r\} \rho_1$$



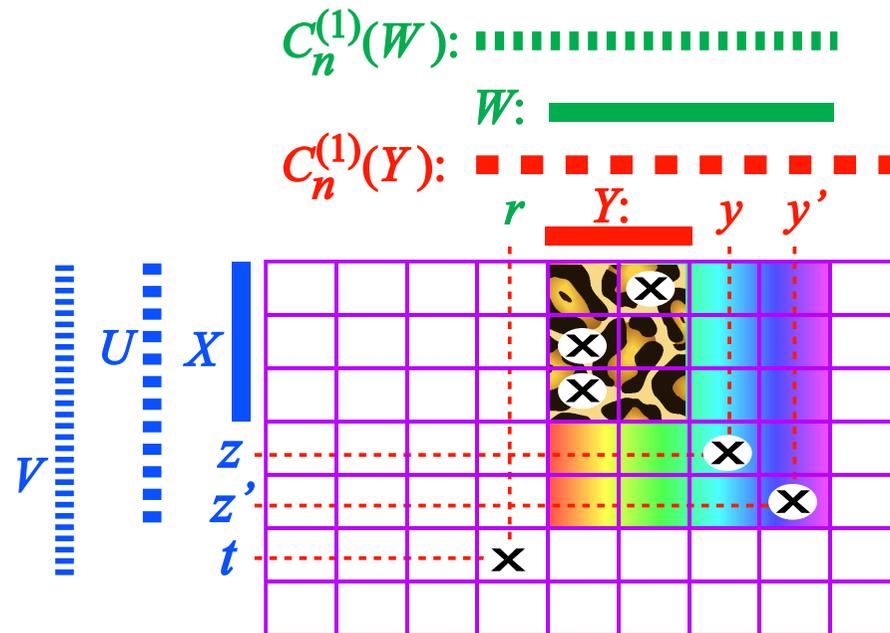
$W \in \psi(U)$ and the formula above for $W = c_{n+1}^{(0)}(U)$ yield: \exists

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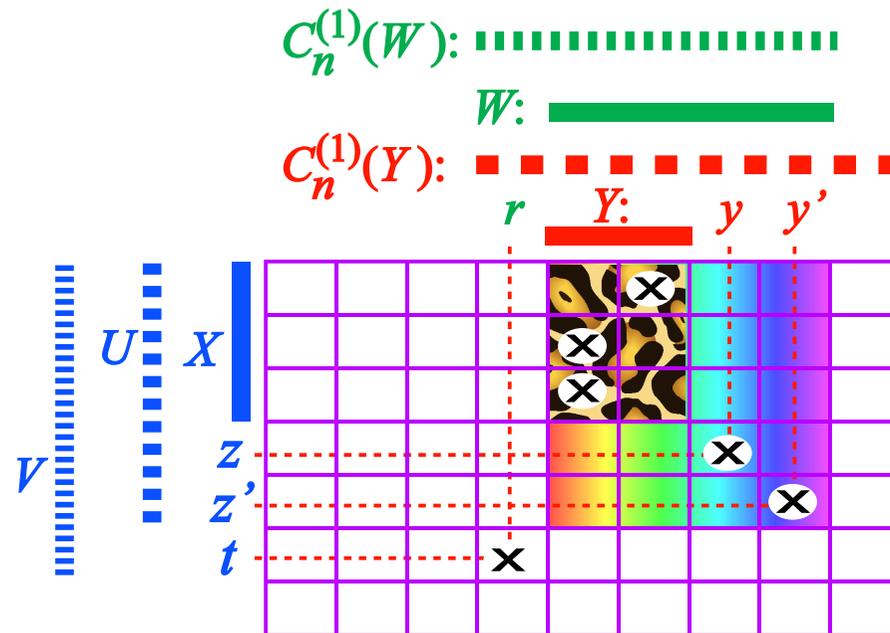
$$r \in c_n^{(1)}(W) = c_n^{(1)}(Y).$$

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$r \in C_n^{(1)}(W) = C_n^{(1)}(Y). \forall Y \implies t \in C_{n+1}^{(0)}(X) = U. \text{ Contradiction.}$

The omitted details are similar (but easier). **Q.e.d.**

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Now let us see an

Example:

The diagram illustrates a grid with a red vertical bar on the left and blue vertical bars between the grid and the label X_0 . The grid has columns labeled B , V , R , and H , and rows labeled a , s , r , f , and m . The cells containing 'X' are highlighted in green.

	B	V	R	H
a	X			X
s	X		X	
r	X			
f	X	X		
m		X	X	X

	<i>B</i>	<i>V</i>	<i>R</i>	<i>H</i>
<i>a</i>	X			X
<i>s</i>	X		X	
<i>r</i>	X			
<i>f</i>	X	X		
<i>m</i>		X	X	X

Objects: **a**pple-eating, **s**nake, **r**obot, **f**ountain, „**m**ixed”

Attributes: **b**alls, **v**ertical, **r**esting time, **h**ead touch.

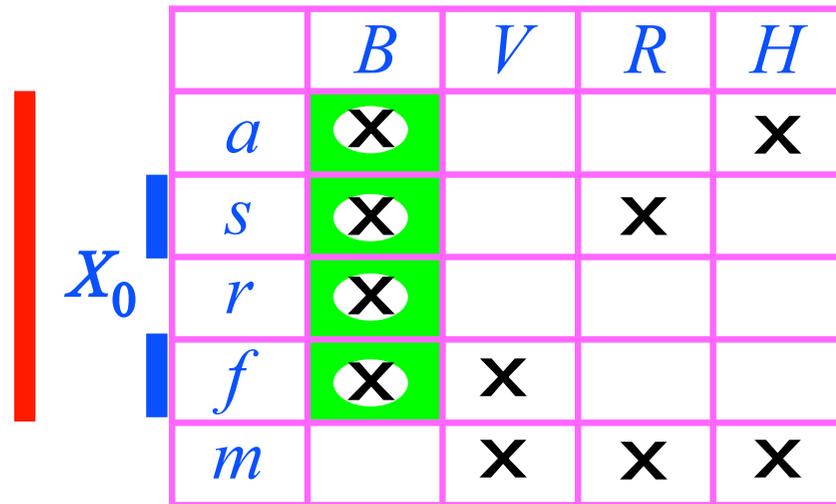
	<i>B</i>	<i>V</i>	<i>R</i>	<i>H</i>
<i>a</i>	X			X
<i>s</i>	X		X	
<i>r</i>	X			
<i>f</i>	X	X		
<i>m</i>		X	X	X

Question: what can we associate with $X = X_0 = \{s, f\}$? Or: if we can perform **S**nake and **f**ountain, what should we learn next?

	B	V	R	H
a	X			X
s	X		X	
r	X			
f	X	X		
m		X	X	X

Compute the G–W closure:

	<i>B</i>	<i>V</i>	<i>R</i>	<i>H</i>
<i>a</i>	X			X
<i>s</i>	X		X	
<i>r</i>	X			
<i>f</i>	X	X		
<i>m</i>		X	X	X



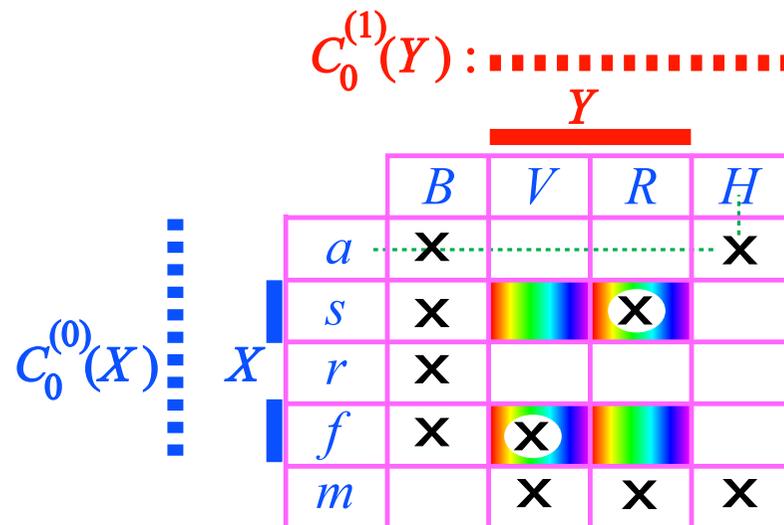
Compute the G–W closure: $C_0^{(0)}(\{s, f\}) = \{s, f, a, r\}$. This says **a**pple-eating or **r**obot. **And we are still hesitating.**

	B	V	R	H
a	X			X
s	X		X	
r	X			
f	X	X		
m		X	X	X

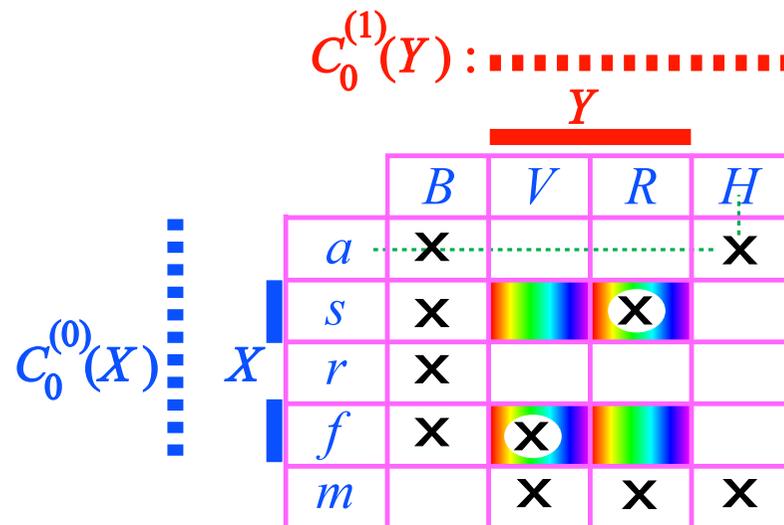
X_0

Compute the G–W closure: $C_0^{(0)}(\{s, f\}) = \{s, f, a, r\}$. This says **a**pple-eating or **r**obot. **And we are still hesitating.** Calculate $C_1^{(0)}(\{s, f\})!$

$$C_{n+1}^{(i)}(X) := C_n^{(i)}(X) \cap \bigcap_{Y \in X \psi_i} \bigcup_{y \in C_n^{(1-i)}(Y)} \{y\} \rho_{1-i}$$

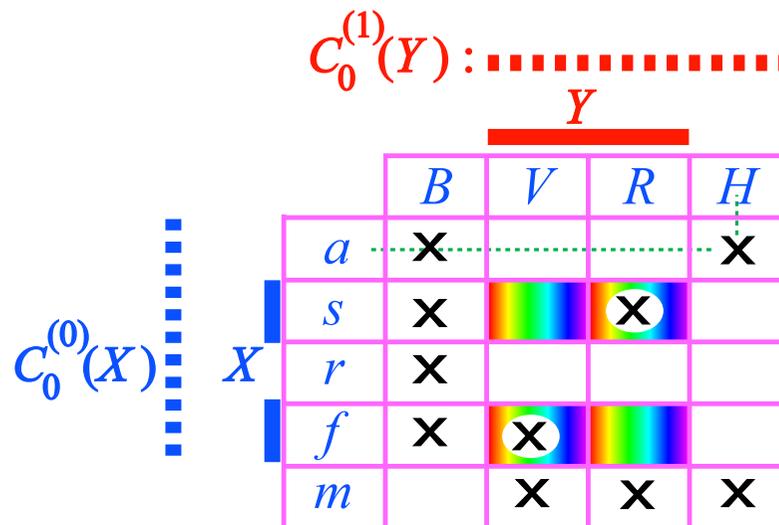


$$C_{n+1}^{(i)}(X) := C_n^{(i)}(X) \cap \bigcap_{Y \in X \psi_i} \bigcup_{y \in C_n^{(1-i)}(Y)} \{y\} \rho_{1-i}$$



Observe:

$$C_{n+1}^{(i)}(X) := C_n^{(i)}(X) \cap \bigcap_{Y \in X \psi_i} \bigcup_{y \in C_n^{(1-i)}(Y)} \{y\} \rho_{1-i}$$



Observe: $a \in C_1^{(0)}(X)$ but $r \notin C_1^{(0)}(X)$.

The final result: $c^{(0)}(X) = c_1^{(0)}(X) = \{s, f\}$,

The final result: $c^{(0)}(X) = c_1^{(0)}(X) = \{s, f, a\}$.

Thus, in the stronger sense, the **a**pple-eating can be associated with $X = \{\mathbf{s}$ nake, **f**ountain} but the **r**obot cannot.

So we should choose the **a**pple-eating.

More serious applications

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For lattices we have:

1. Proposition. *If L is a finite **modular** lattice and we consider the context (J, M, \leq) then \mathcal{C}_0 (the pair of G–W closures) equals \mathcal{C} .*

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The converse is false. \implies **Problem:** Find a better statement.

According to Wille, (J, M, \leq) is interesting.

Since lattices have majority term, the rest of this talk also gives another connection between lattices and \mathcal{C} .

In many cases, \mathcal{C} coincides with $\mathcal{C}^{(0)}$. It is trivial that:

				x				
				x				
b_0	x							
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In this case $\mathcal{C}_0 = \mathcal{C}_n = \mathcal{C}$.

Majority term versus \mathcal{C}

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A ternary term m of an algebra A is called a **majority term** if

$$m(x, x, y) = m(x, y, x) = m(y, x, x) = x$$

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Example: lattices, where $m(x, y, z) = (x \wedge y) \vee (x \wedge z) \vee (y \wedge z)$.

2. Proposition. *Let $A^{(0)}$ and $A^{(1)}$ be algebras with a majority term m , and let $\rho \leq_{sd} A^{(0)} \times A^{(1)}$ be a subdirect product of them. Then for any $i \in \{0, 1\}$ and any $a, b, c \in A^{(i)}$ we have*

$$m(a, b, c) \in \mathcal{C}^{(i)}(\{a, b\}) \cap \mathcal{C}^{(i)}(\{a, c\}) \cap \mathcal{C}^{(i)}(\{b, c\}),$$

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1. Corollary. *If $\rho \leq_{\text{sd}} A^{(0)} \times A^{(1)}$ has a majority term then **all** the triangles of $(A^{(0)}, A^{(1)}, \rho)$ are degenerate.*

Proof of Proposition 2: Let $a_0, b_0, c_0 \in A^{(0)}$, $u = m(a_0, b_0, c_0)$ and $X := \{a_0, b_0\}$. By symmetry, it suffices to show that

$$u \in \mathcal{C}_n^{(0)}(X), \quad \text{for all } n.$$

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let $y \in X\rho_0$ be arbitrary. I.e., $(a_0, y), (b_0, y) \in \rho_0 = \rho$. Since ρ is a subdirect product, $\exists z$ with $(c_0, z) \in \rho$. Then

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shows $u \in \{y\}\rho_1$. This settles the case $n = 0$.

$$c_{n+1}^{(0)}(X) := c_n^{(0)}(X) \cap \bigcap_{Y \in \psi(X)} \bigcup_{y \in c_n^{(1)}(Y)} \{y\} \rho_1$$

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Step $n \rightarrow n + 1$:

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i.e., $u \in \{y\} \rho_1$. Thus, $u \in C_{n+1}^{(0)}(X)$. Q.e.d.

Def. $\alpha \in \text{Con}(A)$ is called
2-uniform if each α -block is 2-element;
k-uniform if each α -block is k-element;
uniform if all α -blocks have the same number of elements.
 A is called **uniform** if all $\alpha \in \text{Con}(A)$ are uniform.

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Grätzer, Quackenbush and Schmidt asked: are finite uniform lattices congruence permutable?

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He also proved that „finiteness” cannot be omitted.

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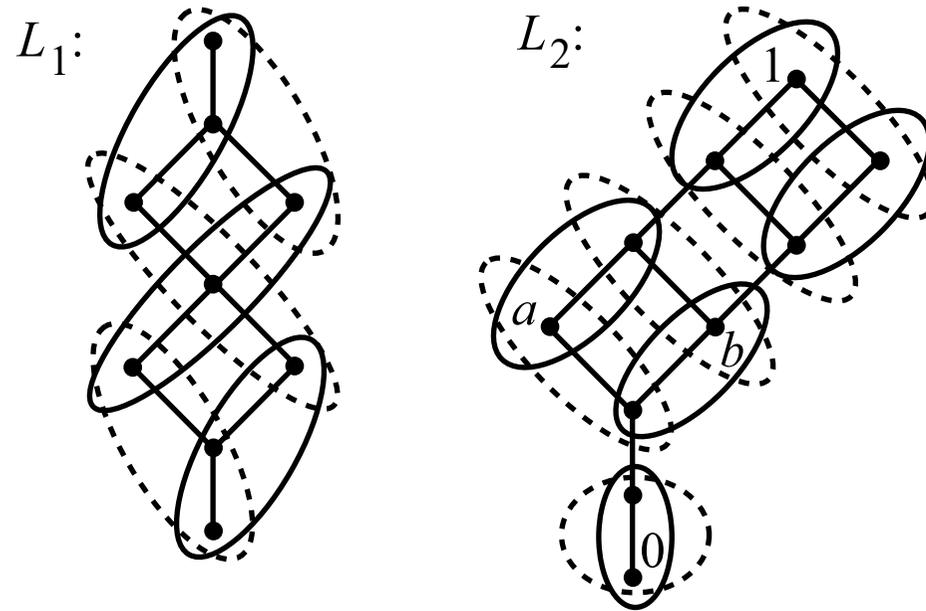
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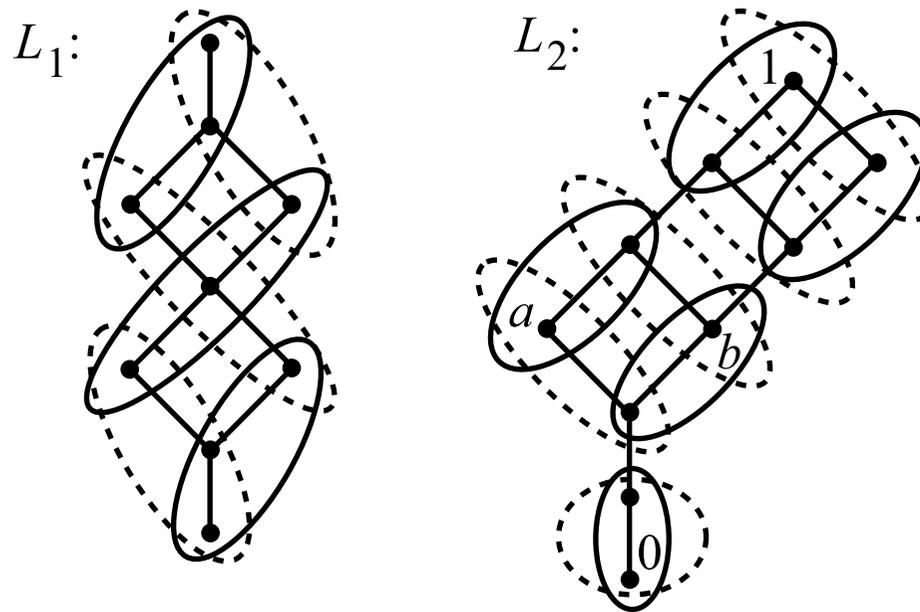
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Remarks: (1) But I will prove this one. (2) „Finiteness” is essential. (3) This thm gives something new even for lattices:



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1. Our main theorem. *Let A be a finite algebra and let V denote the variety generated by A . If*

$$1 \notin \text{typ } V$$

then any two 2-uniform congruences of A permute.

This is a much stronger statement. The proof uses the previous (weaker) theorem.

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Kaarli/1: for $\alpha, \beta \in \text{Con}(A)$ we may assume that $\alpha \vee \beta = 1$. Otherwise take the $\alpha \vee \beta$ blocks. They are subalgebras by idempotency, and it suffices to show that these subalgebras are CP.

Kaarli/2: We may suppose $\alpha \wedge \beta = 0$. Otherwise we go to $A/(\alpha \wedge \beta)$, and 2-uniformity *on nontrivial blocks* is preserved.

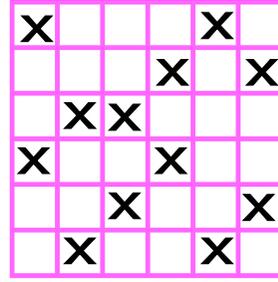
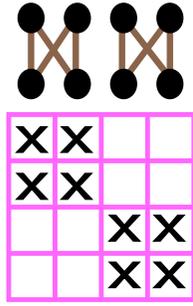
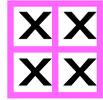
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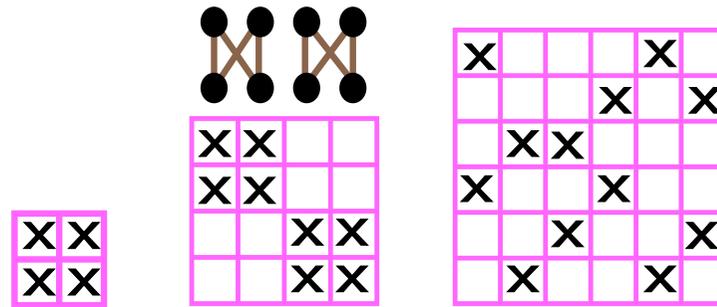
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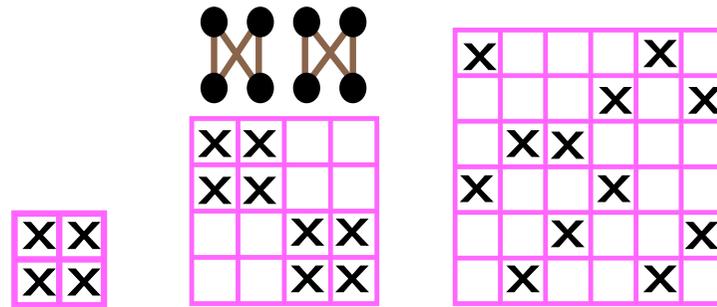
Kaarli/3: $\alpha \wedge \beta = 0$ gives rise to a subdirect product.

These three ideas give easily that only the following proposition needs proving. We call **a context 2-uniform** if each row and each column has exactly two crosses. For example:



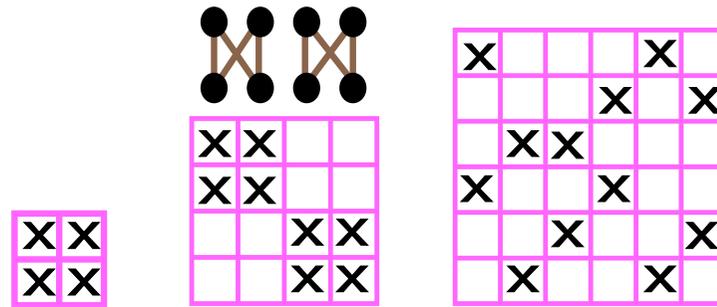


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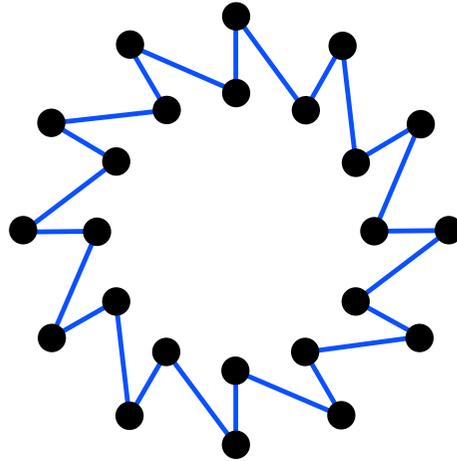
(2) The second one, also given by a bipartite graph, does not correspond to $\alpha \vee \beta = 1$, since the graph **is not connected**. Such contexts are called **decomposable**.

(3) Only the 3rd one has to be treated. This will be done by

3. Proposition. *Each finite 2-uniform indecomposable context $(A^{(0)}, A^{(1)}, \rho)$ with $|A^{(i)}| \geq 3$ contains a nondegenerate triangle.*

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Proof We prove more, namely: **\mathcal{C} is trivial!** Since the context (as a bipartite graph) is 2-regular and connected, it has a Hamiltonian circle. So it can be depicted this way:

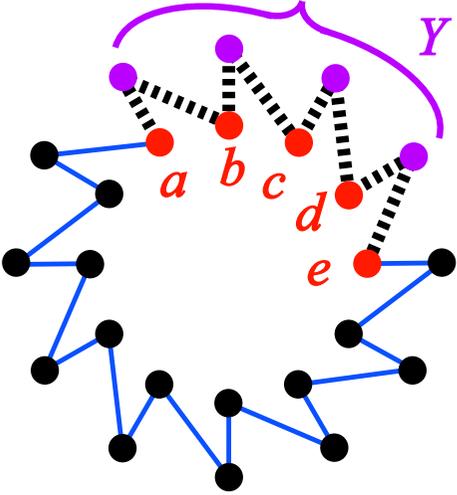


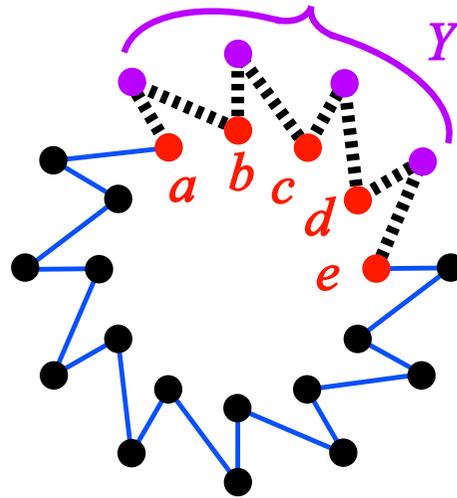
The inner vertices of this „circular saw” constitute $A^{(0)}$ while the outer ones form $A^{(1)}$.

First we show that $\mathcal{C}^{(i)}(X) = X$ for any **arc** in $P(A^{(i)})$. (I.e., \mathcal{C} is trivial on arcs.) Via induction on $|X|$, the length of the arc.

The task is trivial when $|X| < 3$; indeed, then even $\mathcal{C}_0^{(i)}(X) = X$ holds.

Now suppose we have a longer arc, say $X = \{a, b, c, d, e\}$.





Then $Y \in \psi_{\min}(X)$. Ind.hyp. $\implies \mathcal{C}_n^{(1)}(Y) = Y$. Observe that $\{y\}\rho_1 \subseteq X$ for any $y \in Y = \mathcal{C}_n^{(1)}(Y)$. Hence

$$\mathcal{C}_{n+1}^{(0)}(X) \subseteq \bigcap_{Y \in \psi_{\min}(X)} \bigcup_{y \in \mathcal{C}_n^{(1)}(Y)} \{y\}\rho_1 \subseteq X,$$

proving $\mathcal{C}^{(0)}(X) = X$. Thus $\mathcal{C}^{(i)}(X) = X$ for any arc $X \in P(A^{(i)})$.

Finally, each $Y \in P(A^{(i)})$ is an intersection of arcs. \Rightarrow

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Using the fact that $\mathcal{C}^{(i)}$ is a closure operator, it is routine to conclude that $\mathcal{C}^{(i)}(Y) = Y$ for all $Y \in P(A^{(i)})$.

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This proves the proposition and the theorem.

Further open problem: Does any finite k -uniform indecomposable context with $|A^{(i)}| > k$ contain a nondegenerate triangle?

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Further open problem: Does any finite k -uniform indecomposable context with $|A^{(i)}| > k$ contain a nondegenerate triangle?

(Thousands of randomly chosen k -uniform contexts have been tested. An affirmative answer would strengthen Kaarli's result from lattices to algebras with majority term.)

And another open problem: How far can Kaarli's result be generalized? *Majority algebras? CD?*

time?

Gábor Czédli

An association rule ...

Grätzer–Schmidt conf. Budapest 2006 41'

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(In spite of several asymptotic results on the length of the longest circle in the covering graph of the poset

$$(\{X : X \subseteq \{1, 2, \dots, 2k + 1\} \text{ and } |X| \in \{k, k + 1\}\}; \subseteq),$$

it is not known if this graph has a Hamiltonian circle.)

time??

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