

CSPs over structures with slow unlabelled growth

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Constraint Satisfaction Problems

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Alternative formulation

- INPUT: A **primitive positive** sentence φ over \mathfrak{B}
($\varphi \equiv \exists \exists \dots \exists (\bigwedge (\text{atomic}))$)
- QUESTION: Is φ true in \mathfrak{B} ?

Constraint Satisfaction Problems

Finite domain dichotomy

Theorem (Bulatov; Zhuk '17)

If \mathcal{A} is *finite* then $\text{CSP}(\mathcal{A})$ is in **P** or it is **NP**-complete.

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Infinite domain dichotomy

Conjecture (Bodirsky, Pinsker '11)

If \mathfrak{A} is a reduct of a finitely bounded homogeneous structure then $\text{CSP}(\mathfrak{A})$ is in **P** or it is **NP**-complete.

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Remark: all these structures are ω -categorical.

Definition

\mathfrak{A} is ω -categorical if $\text{Aut}(\mathfrak{A})$ has finitely many n -orbits for all $n \in \omega$.

Constraint Satisfaction Problems

Infinite domain dichotomy

CSP dichotomy is solved for

- reducts of $(\mathbb{N}, =)$ (Bodirsky, Kára '08)
- reducts of $(\mathbb{Q}, <)$ (Bodirsky, Kára '09)
- reducts of the homogeneous binary branching C-structure (Bodirsky, Jonsson, Pham '16)
- reducts of the random graph (Bodirsky, Martin, Pinsker, Pongrácz '19)
- reducts of the random poset (Kompatscher, Pham '18)
- reducts of unary ω -categorical structures (Bodirsky, Mottet '18)
- MMSNPs (Bodirsky, Madelaine, Mottet '18)
- reducts the random tournament (Mottet, Pinsker '21)
- first-order expansions of the homogeneous RCC5 structure (Bodirsky, B. '21)
- hereditarily cellular structures (B. '22)
- first-order expansions powers of $(\mathbb{Q}, <)$ (Bodirsky, Jonsson, Martin, Mottet, Semanišínová)
- reducts of random uniform hypergraphs (Mottet, Nagy, Pinsker)

Orbit growth function

Definition

\mathfrak{A} : structure.

- $o_n(\mathfrak{A}) := \#\{n\text{-orbits of } \mathfrak{A}\} = \#\{\text{orbits of } \text{Aut}(\mathfrak{A}) \curvearrowright A^n\}$

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General question:

Given some upper bound $f: \omega \rightarrow \mathbb{R}$, determine all structures with $\ell_n(\mathfrak{A}) < f(n)$ or $u_n(\mathfrak{A}) < f(n)$.

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finite structures + reducts of $(\mathbb{N}, =, c_1, \dots, c_n)$ (Easy.)

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- $u_n < c^n$ for all/some $c > 1$: Topic of today's talk, examples later.

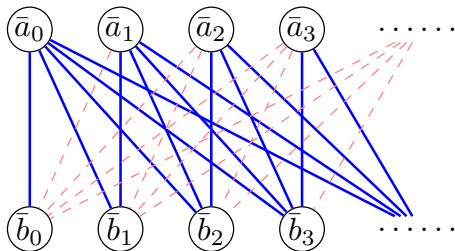
Definition

A formula $\phi(\bar{x}, \bar{y})$ has the **order property** (in \mathfrak{A}) if for all $\exists(\bar{a}_j, \bar{b}_j : j \in \omega)$ such that $\mathfrak{A} \models \phi(\bar{a}_i, \bar{b}_j) \Leftrightarrow i \leq j$. (“ ϕ defines a half-graph”)

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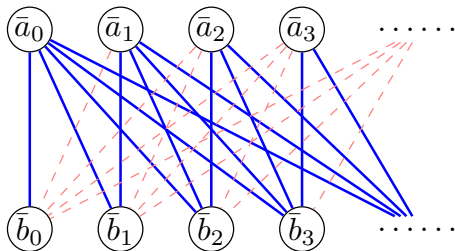


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A structure \mathfrak{A} is **stable** if no formula in \mathfrak{A} has the order property.

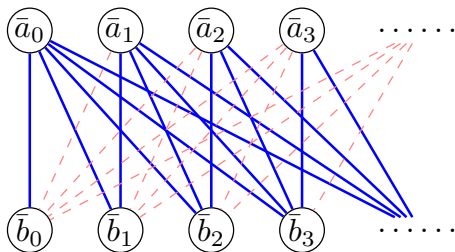


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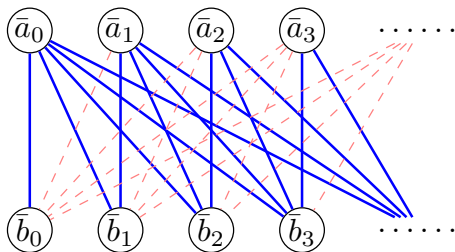
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A structure \mathfrak{A} is **stable** if no formula in \mathfrak{A} has the order property.



- Stable structures: pure set, unary structures, vector spaces, ...
- Unstable structures: $(\mathbb{Q}; <)$, random graph, infinite Boolean algebras, ...

Stable case

Orbit growth dichotomy

Theorem (Braunfeld '22)

\mathfrak{A} : ω -categorical stable structure. Then one of the following holds.

- 1 $\forall c > 1 (u_n(\mathfrak{A}) < c^n)$.
- 2 $\forall c (u_n(\mathfrak{A}) > c^n)$.

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- 1 $\forall c > 1 (u_n(\mathfrak{A}) < c^n)$.
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Moreover the 1 holds if and only if \mathfrak{A} is *hereditarily cellular* (see next slide).

Hereditarily cellular structures

Recursive description

Definition/Theorem (B. '21+Lachlan '92)

\mathfrak{A} is **hereditarily cellular** iff it can be constructed from finite structures by taking

- 1 finite disjoint unions,
- 2 infinite copies,
- 3 first-order reducts.

Hereditarily cellular structures

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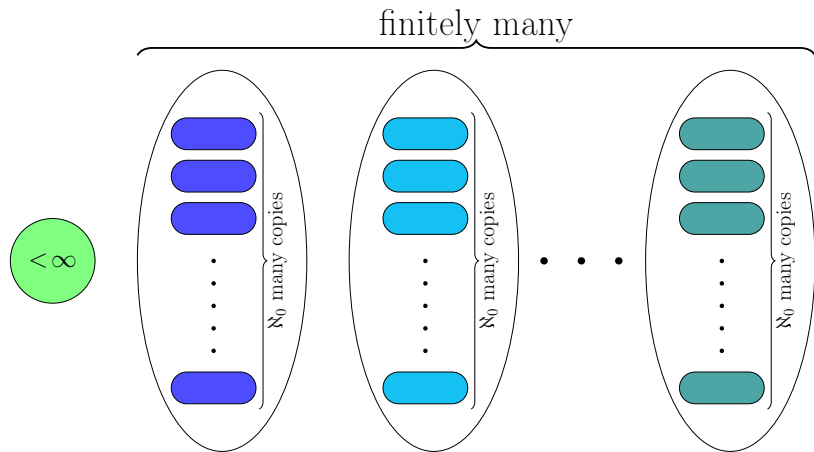
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- **Infinite copies**: we add an equivalence relation E whose equivalence classes are the copies

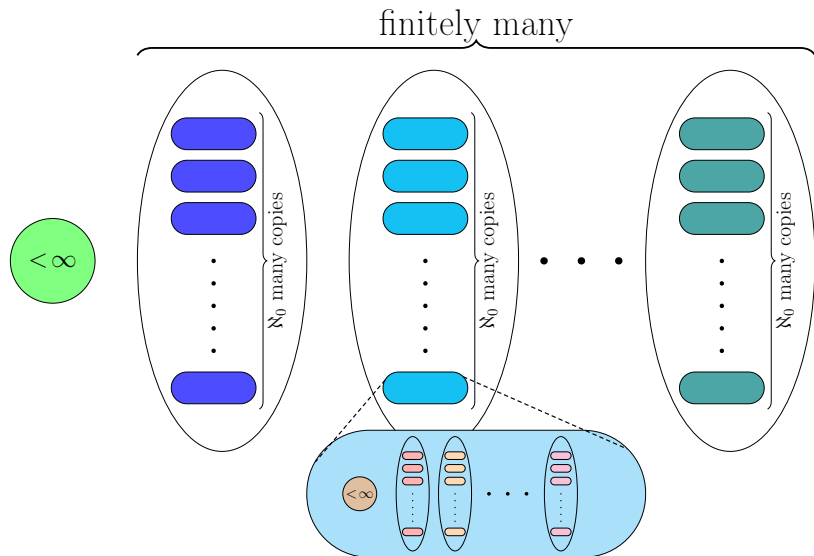
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Illustration



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Getting back to CSPs

Fact (B.)

All hereditarily cellular structures are interdefinable with a finitely bounded homogeneous structure.

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They all fall into the infinite-domain dichotomy conjecture.

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Theorem (B. '21)

*For every hereditarily cellular structure \mathfrak{A} (with a finite relational signature) $\text{CSP}(\mathfrak{A})$ is in **P** or **NP**-complete.*

Hereditarily cellular structures

Proof of the CSP dichotomy

Idea of the proof:

- By general results (Barto, Pinsker) we can assume that \mathfrak{A} has pseudo-Siggers polymorphism:

$$(\alpha \circ f)(x, y, x, z, y, z) = (\beta \circ f)(y, x, z, x, z, y) : \alpha, \beta \in \text{End}(\mathfrak{A})$$

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Remark: this always works for hereditarily cellular structures (heaven).

Structures with subexponential unlabelled growth

Removing stability...

Theorem (B. '22)

$u_n(\mathfrak{A}) < c^n$ for all $c > 1$ iff it can be constructed from finite structures and $(\mathbb{Q}, <)$ by taking

- 1 finite disjoint unions,
- 2 infinite copies,
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Structures with at most $(2 - \varepsilon)^n$ unlabelled growth

Theorem (B. '22)

If $u_n(\mathfrak{A}) < c^n$ for some $c < 2$ iff it can be constructed from finite structures and “some” finite covers of $(\mathbb{Q}, <)$ by taking

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there exists a definable equivalence relation E with finite classes such that $\text{Aut}(\mathfrak{A})/E \simeq \text{Aut}(\mathbb{Q})$.

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Corollary (B. '22)

These structures are all interdefinable with a finitely bounded homogeneous structure.

We are still in scope of the CSP dichotomy conjecture.

Structures with at most $(2 - \varepsilon)^n$ unlabelled growth

Growth rates

Theorem

Assume $u_n(\mathfrak{A}) < c^n$ for some $c < 2$. Then one of the following holds.

- 1 $u_n(\mathfrak{A})$ is slower than exponential.
- 2 $c_1 \gamma_d^n < u_n(\mathfrak{A}) < c_2 (\gamma_d + o(1))^n$ for some $2 \leq d \in \omega$ where γ_d is the largest real root of the polynomial $x^d - x^{d-1} - \dots - 1 = 0$.

$$\gamma_2 \approx 1.618$$

$$\gamma_3 \approx 1.839$$

$$\gamma_d \nearrow 2$$

Structures with at most $(2 - \varepsilon)^n$ unlabelled growth

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Bad news:

general reduction to finite domain CSPs never works in the unstable case (hell)!

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*Let \mathfrak{A} be a first-order expansion of $(\mathbb{Q}, <)$. Then if $\text{CSP}(\mathfrak{A})$ is not **NP**-complete then $\text{Pol}(\mathfrak{A})$ contains pp, ll, dual-pp, dual-ll or a constant operation.*

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- ω -categorical unary expansions of $(\mathbb{Q}, <)$?

Further questions

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More examples:

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- homogeneous C - and D -structures
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Conjecture (Braunfeld)

The class $\{\mathfrak{A} : u_n(\mathfrak{A}) < c^n \text{ for some } c\}$ is closed under taking ω -categorical unary expansions.

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More examples:

- finite covers of $(\mathbb{Q}, <)$
- $S(2)$, the *local order*, $u_n \sim 2^{n-1}/n$
- homogeneous C - and D -structures
- we can take finite covers and expansions with *random partitions*

Conjecture (Braunfeld)

The class $\{\mathfrak{A} : u_n(\mathfrak{A}) < c^n \text{ for some } c\}$ is closed under taking ω -categorical unary expansions.

What are the possible growth rates of structures with $u_n < c^n$?

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Conjecture

If $u_n(\mathfrak{A}) < c^n$ then \mathfrak{A} is finitely homogenizable.