# Lattice-induced threshold functions and Boolean functions

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## Threshold functions

A classical **threshold function** is a Boolean function  $f: \{0,1\}^n \to \{0,1\}$  such that there exist real numbers  $w_1, \ldots, w_n, t$ , fulfilling

$$f(x_1, \ldots, x_n) = 1$$
 if and only if  $\sum_{i=1}^n w_i \cdot x_i \ge t$ ,

where  $w_i$  is called **weight** of  $x_i$ , for i = 1, 2, ..., n and t is a constant called the **threshold value**.

modeling neurons
political decisions
electrical engineering
artifical intelligence
game theory

## Threshold functions

- combinatorics (their number!!!)
- computer science

#### **IN ALGEBRA**

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tolerance relation (B. Bodi)
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fundamental ideal of a groupring (B. Bodi)

generalized clones (constraints) (S. Foldes, L. Hellerstein, M. Couceiro) no superposition, not clone

invariance group

coalition lattice (conjecture)

# Monotonicity and thresholdness

It is easy to see that threshold functions with positive weights and a threshold value are isotone.

However, an isotone Boolean function is not necessarily threshold, e.g.  $f = x \cdot y \lor w \cdot z$  is isotone, but not a threshold function.

## Threshold functions

 $f = x \cdot y \lor w \cdot z$  is isotone, but not a threshold function because its invariance group is

D8 =

$$\{(), (1324), (12)(34), (1423), (12), (34), (12)(34), (13)(24), (14), (23)\}$$

## Theorem (1994.)

For every n-ary threshold function f there exists a partition  $C_f$  of the set of variables X such that the invariance group G of f consists of exactly those permuations of  $S_X$  which preserve each block of  $C_f$ .

I.e. the invariance groups of threshold functions are of the following form: direct product of symmetric groups.

## Lattice-induced threshold functions

Let L be a complete lattice in which the bottom and the top are (also) denoted by 0 and 1 respectively; however, it is clear from the context whether 0 (1) is a component in some  $(x_1, \ldots, x_n) \in \{0, 1\}^n$ , or it is from L.

For  $x \in \{0,1\}$ , and  $w \in L$ , we define a mapping  $L \times \{0,1\}$  into L denoted by " $\cdot$ ", as follows:

$$w \cdot x := \begin{cases} w, & \text{if } x = 1 \\ 0, & \text{if } x = 0. \end{cases}$$
 (1)

A function  $f:\{0,1\}^n \to \{0,1\}$  is a **lattice-induced threshold function**, if there is a complete lattice L and  $w_1,\ldots,w_n,t\in L$ , such that

$$f(x_1,\ldots,x_n)=1$$
 if and only if  $\bigvee_{i=1}^n(w_i\cdot x_i)\geq t.$  (2)

## Lattice-induced threshold functions

#### **Proposition**

Every lattice-induced threshold function is isotone.

#### Theorem

Every isotone Boolean function is a lattice-induced threshold function.

#### Remark

The corresponding lattice in each case can be the free distributive lattice with n generators.

# Sketch of the proof of the Theorem

We prove that for every  $n \in \mathbb{N}$ , there is a lattice L such that every isotone Boolean function is a lattice induced threshold function over L. Let  $n \in \mathbb{N}$ .

We take L to be a free distributive lattice with n generators  $w_1, w_2, \ldots, w_n$ .

Recall that every element in a free distributive lattice can be uniquely represented in a "conjunctive normal form" by means of generators (i.e., every element is a meet of elements of the type  $\bigvee_{i \in J} w_i$ , where  $J \subseteq \{1, \ldots, n\}$ .)

For  $x,y\in L$ , if  $x=\bigwedge_{k=1}^p\bigvee_{j\in I_k}w_j$  and  $y=\bigwedge_{k=1}^l\bigvee_{s\in J_k}w_s$ ,  $x\leq y$  if and only if for every  $u\in\{1,\ldots,l\}$  there is  $k\in\{1,\ldots,p\}$  such that  $I_k\subseteq J_u$ . (\*)

# Sketch of the proof of the Theorem

Let  $f: \{0,1\}^n \to \{0,1\}$  be an isotone Boolean function. Let F be the corresponding order semi-filter on  $\{0,1\}^n$ .

Further, let  $m_1, \ldots, m_p$  be minimal elements of this semi-filter. Let  $I_1, \ldots, I_p$  be subsets of  $\{1, 2, \ldots, n\}$ , i.e., sets of indices, such that  $i \in I_k$  if and only if i - th coordinate of  $m_k$  is equal to 1.

For the threshold  $t \in L$  associated to the given function f we take

$$t = \bigwedge_{k=1}^{p} \bigvee_{j \in I_k} w_j.$$

Now, it is straightforward to show that

$$f(x_1,\ldots,x_n)=1$$
 if and only if  $\bigvee_{i=1}^n(w_i\cdot x_i)\geq t.$  (3)

## Lattice-valued Boolean functions, cuts

A function  $f: \{0,1\}^n \to L$ , where L is a complete lattice, is called a **lattice valued** (L-valued) **Boolean function**.

For  $f:\{0,1\}^n \to L$  and  $p \in L$ , a cut set (cut)  $f_p$  is a subset of  $\{0,1\}^n$ :

$$f_p = \{x \in \{0,1\}^n \mid f(x) \ge p\}.$$

In other words, a p-cut of  $\mu: B \to L$  is the inverse image of the principal filter  $\uparrow p$ , generated by  $p \in L$ :

$$\mu_{p} = \mu^{-1}(\uparrow p). \tag{4}$$

It is obvious that for  $p, q \in L$ ,

from  $p \leq q$  it follows that  $\mu_q \subseteq \mu_p$ .

## Lattice-valued up-sets

An *L*-valued Boolean function  $\mu: B \to L$  is called a **lattice valued** (*L*-valued) up-set, if from  $x \le y$  it follows that  $\mu(x) \le \mu(y)$ .

#### Lemma

Let B be a Boolean lattice and  $\mu: B \to L$  an L-valued Boolean function. Then  $\mu$  is an L-valued up-set on B if and only if all the cuts of  $\mu$  are up-sets (order-filters, semi-filters) on B.

# Representation of lattice-valued up-sets by cuts

Let  $B=(\{0,1\}^n,\leq)$ ,  $n\in\mathbb{N}$ ,  $L_D$  a free distributive lattice with n generators  $w_1,\ldots,w_n$  and  $\overline{\beta}:B\to L_D$ , an  $L_D$ -valued function on B defined in the following way: for  $x=(x_1,\ldots,x_n)\in B$ 

$$\overline{\beta}(x) = \bigvee_{i=1}^{n} (w_i \cdot x_i), \tag{5}$$

where the function " $\cdot$ " is defined by (1). By the definition,  $\overline{\beta}$  is uniquely (up to a permutation of generators  $w_i$ ) determined by a finite Boolean lattice  $B = (\{0,1\}^n, \leq)$ , i.e., by a positive integer n.

#### **Observations**

The  $L_D$ -valued function  $\overline{\beta}$  defined by (5) is an  $L_D$ -valued up-set on B. Every cut of  $\overline{\beta}$  is an up-set of a finite Boolean lattice  $B = (\{0,1\}^n, \leq), n \in \mathbb{N}$ .

# Representation of lattice-valued up-sets by cuts

#### **Theorem**

Every up-set of a finite Boolean lattice  $B = (\{0,1\}^n, \leq), n \in \mathbb{N}$ , is a cut of  $\overline{\beta}$ .

## **Corollary**

The collection of cuts of every L-valued up-set on B (for any L) is contained in the collection of cuts of  $\overline{\beta}$ .

## Linear combinations

Let  $B=(\{0,1\}^n,\leq)$  be a Boolean lattice, L a complete lattice,  $x=(x_1,\ldots,x_n)\in B$  and  $w_1,\ldots,w_n\in L$ . Further, let the binary function "·" which maps  $L\times\{0,1\}$  into L be defined by (1). Then we call the term

$$\bigvee_{i=1}^{n} (w_i \cdot x_i), \tag{6}$$

a **linear combination** of elements  $w_1, \ldots, w_n$  from L.

Observe also that in the case of formula (5), the corresponding  $L_D$ -valued function is  $\overline{\beta}$  and the following is obviously true: the closure system consisting of all up-sets on B is the collection of cuts of  $\overline{\beta}$ .

## Generalization

What about taking an arbitrary lattice L instead of  $L_D$ ? Or, starting with a closure system  $\mathcal F$  consisting of some up-sets on  $B=(\{0,1\}^n,\leq)$ , and we try to find a lattice L and  $w_1,\ldots,w_n\in L$ , such that the family of cuts of the function

$$\bigvee_{i=1}^{n} (w_i \cdot x_i), \tag{7}$$

over this lattice (a linear combination of elements from L) coincides with  $\mathcal{F}$ .

The answer to the above problem is not generally positive, as shown by the following example.

## Example

Let  $B = (\{0,1\}^2, \leq)$  be the four element Boolean lattice and

$$\mathcal{F} = \{\{(1,1)\}, \{(1,1), (1,0)\}, \{(1,1), (1,0), (0,1)\}, \{(1,1), (1,0), (0,1), (0,0)\}\}$$

a closure system consisting of some up-sets on B.

We show that there is no lattice L, hence neither there is an L-valued function  $\nu: B \to L$ , such that there are  $w_1, w_2 \in L$  fulfilling that for all  $x_1, x_2 \in \{0, 1\}$ 

$$\nu(x_1,x_2)=(w_1\cdot x_1)\vee(w_2\cdot x_2)$$

and that the collections of cuts of  $\nu$  is  $\mathcal{F}$ .

## Example

Indeed, suppose that there is a lattice L and elements  $w_1, w_2 \in L$ , such that  $\nu(x_1, x_2) = (w_1 \cdot x_1) \vee (w_2 \cdot x_2)$ , for all  $x_1, x_2 \in \{0, 1\}$ . Then,  $\nu(0,0) = 0 \in L_1$ ,  $\nu(0,1) = w_2$ ,  $\nu(1,0) = w_1$  and  $\nu(1,1) = w_1 \vee w_2$ . Now, since the cuts of  $\nu$  are supposed to be elements from  $\mathcal{F}$ , and cuts are up-sets in B, we have that  $\nu_{w_1 \vee w_2} = \{(1,1)\}$ , and  $w_1 \vee w_2$  would be the top element of the lattice L: otherwise the empty set would be a cut of this lattice valued function.

**Lemma** Let  $\mu: B \to L$  be a lattice valued up-set, such that its collection of cuts is  $\mathcal{F}$ . If  $\uparrow a \in \mathcal{F}$  and  $\mu(a) = p$ , then  $\mu_p = \uparrow a$ .

Now  $\nu_{w_1} = \{(1,1),(1,0)\}, \ \nu_{w_2} = \{(1,1),(1,0),(0,1)\}$  and  $\nu_0 = \{(1,1),(1,0),(0,1),(0,0)\}.$  Since  $(1,0) \in \nu_{w_2}$ , we have that  $\nu(1,0) \geq w_2$ , i.e.,  $w_1 \geq w_2$ . Hence,  $w_1 \vee w_2 = w_1$ , which contradicts the assumption that  $\nu_{w_1\vee w_2} \neq \nu_{w_1}$ .

Hence, the up-sets from the collection  $\mathcal{F}$  cannot be represented as cuts of an L-valued function in the form (7).

# First problem

Find necessary and sufficient conditions under which a lattice valued up-set  $\mu: B \to L$  on a finite Boolean lattice  $B = (\{0,1\}^n, \leq)$  can be represented by the linear combination

$$\mu(x) = \bigvee_{i=1}^{n} (w_i \cdot x_i)$$

over  $L(x = (x_1, ..., x_n) \in \{0, 1\}^n, w_1, ..., w_n \in L$ .

## **Definition**

Starting with finite lattices M and L with the bottom elements  $0_M$  and  $0_L$  respectively, we say that a mapping  $\mu: M \to L$  is a  $0-\vee$ -homomorphism, if for all  $x,y \in M$ 

$$\mu(x \lor y) = \mu(x) \lor \mu(y)$$
 and  $\mu(0_M) = 0_L$ .

In particular, if  $\mu$  maps a Boolean lattice  $B=\{0,1\}^n$  into L, the condition that  $\mu$  is a 0–V–homomorphism from B to L is equivalent with the following two conditions (observe that B is finite): for every collection A of some atoms in B

(i) 
$$\mu(\bigvee A) = \bigvee \mu(A)$$
 and (ii)  $\mu(0, \dots, 0) = 0$ . (8)

## Characterization

Let  $B=(\{0,1\}^n,\leq)$  be a finite Boolean lattice and L an arbitrary complete lattice. Then an L-valued Boolean function  $\mu:\{0,1\}^n\to L$  can be represented in the form

$$\mu(x) = \bigvee_{i=1}^{n} (w_i \cdot x_i)$$

for some elements  $w_1, \ldots, w_n \in L$  if and only if  $\mu$  as a mapping from B to L is a 0– $\vee$ –homomorphism.

## Next problem

Next we analyze the same problem (representability by linear combination), for closure systems of some up-sets on  $B=(\{0,1\}^n,\leq)$  .

If  $\mathcal{F}$  is a closure system consisting of some up-sets on  $B=(\{0,1\}^n,\leq)$ , then for  $x\in P$ , we define

$$\overline{x} = \bigcap \{ f \in \mathcal{F} \mid x \in f \}. \tag{9}$$

**Proposition** Let  $\mathcal{F}$  be a closure system of some up-sets on B. If  $\mathcal{F}$  is a family of cuts of an L-valued up-set  $\mu$  on B represented by a linear combination over L, then the following holds: for all  $x, y \in B$ 

from 
$$\overline{x} \subseteq \overline{y}$$
 it follows that  $\overline{x \vee y} = \overline{x}$ . (10)

# Proof of this Proposition

#### Lemma

Let  $\mathcal{F}$  be a closure system consisting of some up-sets on a poset  $(P, \leq)$ . For  $x \in P$ , denote

$$\overline{x} = \bigcap \{ f \in \mathcal{F} \mid x \in f \}. \tag{11}$$

Then, for all  $x, y, z \in P$ , the following is true:

- a)  $x \le y$  implies  $\overline{y} \subseteq \overline{x}$ .
- b)  $x \in \overline{x}$ .
- c)  $\uparrow x \subseteq \overline{x}$ .
- d) If  $z \in \overline{x}$  then  $\overline{z} \subseteq \overline{x}$ .

# Proof of this Proposition

We introduce the mapping  $\widehat{\mu}: B \to \mu_L$  by the construction by

$$\widehat{\mu}(x) := \bigcap \{ \mu_p \in \mu_L \mid x \in \mu_p \}. \tag{12}$$

We say that the lattice valued function  $\widehat{\mu}$  is the **canonical representation** of  $\mu$ .

#### **Proposition**

If  $\mu: B \to L$  is an L-valued function on B and  $\mu(a) = \mu(b) \vee \mu(c)$  for some  $a, b, c \in B$ , then also for the canonical representation  $\widehat{\mu}$  of  $\mu$ ,  $\widehat{\mu}(a) = \widehat{\mu}(b) \vee \widehat{\mu}(c)$  analogously holds.

### Remark

The opposite implication to the one in this Proposition does not hold in general. Indeed, let  $B = \{a, b, c, d\}$ , and let L be the lattice given in Figure 1.

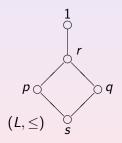


Figure 1

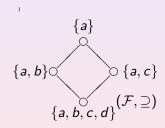


Figure 2

#### Remark

We define an *L*-valued function  $\mu : B \to L$  as follows:

$$\mu = \left(\begin{array}{ccc} a & b & c & d \\ 1 & p & q & s \end{array}\right).$$

The cuts of  $\mu$  are:

 $\mu_L = \{\mu_1 = \mu_r = \{a\}, \mu_p = \{a,b\}, \mu_q = \{a,c\}, \ \mu_s = \{a,b,c,d\}\}.$  The lattice  $(\mu_L,\supseteq)$  is depicted in Figure 2. The canonical representation of this lattice valued function is  $\widehat{\mu}: B \to \mu_L$  and it is given by

$$\widehat{\mu} = \left(\begin{array}{ccc} a & b & c & d \\ \{a\} & \{a,b\} & \{a,c\} & \{a,b,c,d\} \end{array}\right).$$

Now, observe that  $\widehat{\mu}(a) = \widehat{\mu}(b) \vee \widehat{\mu}(c)$ . However, it is not true that  $\mu(a) = \mu(b) \vee \mu(c)$ .

## Example

Let  $B = (\{0,1\}^2, \leq)$  be the four element Boolean lattice and

$$\mathcal{F} = \{\{(1,1)\}, \{(1,1), (1,0)\}, \{(1,1), (1,0), (0,1)\}, \{(1,1), (1,0), (0,1), (0,0)\}, \{(1,1), (1,0), (0,1)\}, \{(1,1), (1,0), (1,0)\}, \{(1$$

a closure system consisting of some up-sets on B.

We already proved that this family is not the collection of cuts for a lattice valued function representable by a linear combination. If we define a mapping from B to  $\mathcal F$  by

$$\overline{x} = \bigcap \{ f \in \mathcal{F} \mid x \in f \},\tag{13}$$

then the condition from Proposition that for all  $x, y \in B$ 

from 
$$\overline{x} \subseteq \overline{y}$$
 it follows that  $\overline{x \vee y} = \overline{x}$ . (14)

is not satisfied.

#### Theorem

Let  $\mathcal{F}$  be a closure system of some up-sets on a Boolean algebra B and for  $x \in B$ , define  $\overline{x}$  by (13):

$$\overline{x} = \bigcap \{ f \in \mathcal{F} \mid x \in f \}.$$

The following conditions are equivalent:

(i) for all  $x, y \in B$ 

from 
$$\overline{x} \subseteq \overline{y}$$
 it follows that  $\overline{x \vee y} = \overline{x}$ .

- (ii) for all  $x, y \in B$ ,  $\overline{x \vee y} = \overline{x} \cap \overline{y}$ .
- (iii) There is a lattice L such that  $\mathcal{F}$  is a family of cuts of an L-valued up-set on B which can be represented as a linear combination over L.

## Last problem

Given a lattice valued up-set  $\mu: B \to L$  on a finite Boolean lattice  $B = \{0,1\}^n$ , find a lattice  $L_1$  and a lattice valued function  $\nu: B \to L_1$  defined by the formula

$$\nu(x) = \bigvee_{i=1}^{n} (w_i \cdot x_i)$$

where  $w_1, \ldots, w_n \in L_1$ , such that the collections of cuts of  $\mu$  and  $\nu$  coincide.

## Last problem solution

#### **Corollary**

For a lattice valued up-set  $\mu: B \to L$  on a finite Boolean lattice  $B = \{0,1\}^n$ , there is a lattice  $L_1$  and a lattice valued function  $\nu: B \to L_1$  defined by the formula

$$\nu(x) = \bigvee_{i=1}^{n} (w_i \cdot x_i)$$

such that the collections of cuts of  $\mu$  and  $\nu$  coincide if and only if  $\overline{x \vee y} = \overline{x} \cap \overline{y}$  for  $x, y \in B$ , where the operator  $\overline{\ }$  is defined by cuts of  $\mu$ : define  $\overline{x}$  by (13):

$$\overline{x} = \bigcap \{ f \in \mathcal{F} \mid x \in f \}.$$

## Lattice-induced threshold functions and Boolean functions

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az Európai Szociális Alap társfinanszírozásával valósul meg.



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