

ALGORITHMIC COMPLEXITY
and UNIVERSAL ALGEBRA

2007. 7. 16. – 7. 20.

Szeged, Hungary

**Some Remarks on
Minimal Clones**

Hajime MACHIDA

(Tokyo)

Joint Work with

Michael PINSKER

(Wien)

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Some Remarks on
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— A ‘Minimal’ Adventure to
the World of Béla Csákány —

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§1 Where to Start

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◇ Béla Csákány ◇

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In 1983, B. Csákány published the following paper :

B. Csákány, “All minimal clones on the three element set”, *Acta Cybernet.*, Vol. 6, 227-238, 1983

where he determined all minimal clones on a three-element set.

Number of minimal clones = 84

§0 Just in Case ... Definitions

clone

$C (\subseteq \mathcal{O}_k) : \text{clone on } E_k$

\iff

(i) $C : \text{closed under composition}$

(ii) $C \supseteq \mathcal{J}_k$

$\mathcal{L}_k : \text{lattice of all clones on } E_k$

minimal clone

$C (\in \mathcal{L}_k) : \text{minimal clone}$

\iff

(i) $C \neq \mathcal{J}_k$

(ii) **for** $\forall C' \in \mathcal{L}_k$,
 $\mathcal{J}_k \subset C' \subseteq C \implies C' = C$

§1 Where to Start ... Continued

♡ Generators of Minimal Clones on a Three-Element Set ♡

(B. Csákány, 1983)

Each of the following operations generates mutually distinct minimal clone.

(II) Binary idempotent operations

$$b_0 = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 1 & 0 \\ \hline 0 & 0 & 2 \\ \hline \end{array} \quad b_{364} = \begin{array}{|c|c|c|} \hline 0 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 1 & 2 \\ \hline \end{array} \quad b_{728} = \begin{array}{|c|c|c|} \hline 0 & 2 & 2 \\ \hline 2 & 1 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_8 = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 1 & 0 \\ \hline 2 & 2 & 2 \\ \hline \end{array} \quad b_{368} = \begin{array}{|c|c|c|} \hline 0 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array} \quad b_{80} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 2 & 1 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{36} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 1 \\ \hline 0 & 0 & 2 \\ \hline \end{array} \quad b_{40} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 1 \\ \hline 1 & 1 & 2 \\ \hline \end{array} \quad b_{692} = \begin{array}{|c|c|c|} \hline 0 & 2 & 2 \\ \hline 1 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{10} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 1 & 1 \\ \hline 0 & 1 & 2 \\ \hline \end{array} \quad b_{280} = \begin{array}{|c|c|c|} \hline 0 & 1 & 0 \\ \hline 1 & 1 & 1 \\ \hline 0 & 1 & 2 \\ \hline \end{array} \quad b_{458} = \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 1 & 1 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{20} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 1 & 2 \\ \hline 0 & 2 & 2 \\ \hline \end{array} \quad b_{448} = \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 1 & 1 & 1 \\ \hline 2 & 1 & 2 \\ \hline \end{array} \quad b_{188} = \begin{array}{|c|c|c|} \hline 0 & 0 & 2 \\ \hline 0 & 1 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{11} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 2 & 2 \end{bmatrix} \quad b_{286} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 2 & 1 & 2 \end{bmatrix} \quad b_{215} = \begin{bmatrix} 0 & 0 & 2 \\ 1 & 1 & 2 \\ 2 & 2 & 2 \end{bmatrix}$$

$$b_{16} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 1 \\ 2 & 1 & 2 \end{bmatrix} \quad b_{281} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 2 & 2 \end{bmatrix} \quad b_{296} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 2 \\ 2 & 2 & 2 \end{bmatrix}$$

$$b_{47} = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 2 \\ 0 & 2 & 2 \end{bmatrix} \quad b_{205} = \begin{bmatrix} 0 & 0 & 2 \\ 1 & 1 & 1 \\ 2 & 1 & 2 \end{bmatrix} \quad b_{179} = \begin{bmatrix} 0 & 0 & 2 \\ 0 & 1 & 1 \\ 2 & 2 & 2 \end{bmatrix}$$

$$b_{17} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 1 \\ 2 & 2 & 2 \end{bmatrix} \quad b_{287} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 2 & 2 & 2 \end{bmatrix} \quad b_{53} = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 2 \\ 2 & 2 & 2 \end{bmatrix}$$

$$b_{38} = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 2 & 2 \end{bmatrix} \quad b_{43} = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 1 \\ 2 & 1 & 2 \end{bmatrix} \quad b_{206} = \begin{bmatrix} 0 & 0 & 2 \\ 1 & 1 & 1 \\ 2 & 2 & 2 \end{bmatrix}$$

$$b_{26} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 2 \\ 2 & 2 & 2 \end{bmatrix} \quad b_{449} = \begin{bmatrix} 0 & 1 & 2 \\ 1 & 1 & 1 \\ 2 & 2 & 2 \end{bmatrix} \quad b_{37} = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 2 \end{bmatrix}$$

$$b_{33} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 0 \\ \hline 2 & 0 & 2 \\ \hline \end{array} \quad b_{122} = \begin{array}{|c|c|c|} \hline 0 & 0 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 2 & 2 \\ \hline \end{array} \quad b_{557} = \begin{array}{|c|c|c|} \hline 0 & 2 & 0 \\ \hline 2 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{35} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 0 \\ \hline 2 & 2 & 2 \\ \hline \end{array} \quad b_{125} = \begin{array}{|c|c|c|} \hline 0 & 0 & 1 \\ \hline 1 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array} \quad b_{71} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 2 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{42} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 1 \\ \hline 2 & 0 & 2 \\ \hline \end{array} \quad b_{41} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 1 \\ \hline 1 & 2 & 2 \\ \hline \end{array} \quad b_{530} = \begin{array}{|c|c|c|} \hline 0 & 2 & 0 \\ \hline 1 & 1 & 1 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{68} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 2 & 1 & 1 \\ \hline 1 & 2 & 2 \\ \hline \end{array} \quad b_{528} = \begin{array}{|c|c|c|} \hline 0 & 2 & 0 \\ \hline 1 & 1 & 1 \\ \hline 2 & 0 & 2 \\ \hline \end{array} \quad b_{116} = \begin{array}{|c|c|c|} \hline 0 & 0 & 1 \\ \hline 1 & 1 & 0 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$b_{178} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 1 & 1 & 0 \\ \hline 2 & 0 & 2 \\ \hline \end{array} \quad b_{290} = \begin{array}{|c|c|c|} \hline 0 & 0 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 2 & 2 \\ \hline \end{array}$$

$$b_{624} = \begin{array}{|c|} \hline 0 \ 2 \ 1 \\ \hline 2 \ 1 \ 0 \\ \hline 1 \ 0 \ 2 \\ \hline \end{array}$$

(III) Ternary majority operations

$$m_0 = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 0 \\ \hline 0\ 0\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 0 \\ \hline 1\ 1\ 1 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 0\ 2 \\ \hline 0\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{364} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 1 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 1 \\ \hline 1\ 1\ 1 \\ \hline 1\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 1\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{728} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 2 \\ \hline 0\ 2\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 1\ 1\ 1 \\ \hline 2\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 2\ 2 \\ \hline 2\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{109} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 0 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 1 \\ \hline 1\ 1\ 1 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 0\ 2 \\ \hline 1\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{473} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 1 \\ \hline 0\ 2\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 1\ 1\ 1 \\ \hline 1\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 2\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{510} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 2 \\ \hline 0\ 0\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 0 \\ \hline 1\ 1\ 1 \\ \hline 2\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 2\ 2 \\ \hline 0\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

$$m_{624} = \begin{array}{|c|} \hline 0\ 0\ 0 \\ \hline 0\ 1\ 2 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 1\ 1\ 1 \\ \hline 0\ 1\ 2 \\ \hline \end{array} \quad \begin{array}{|c|} \hline 0\ 1\ 2 \\ \hline 0\ 1\ 2 \\ \hline 2\ 2\ 2 \\ \hline \end{array}$$

(IV) Ternary semiprojections

$$s_0 = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 0 \\ \hline 1 & 1 & 1 \\ \hline 0 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 0 & 2 \\ \hline 0 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{364} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 1 \\ \hline 0 & 1 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 1 & 2 \\ \hline 1 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{728} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 2 \\ \hline 0 & 2 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 2 \\ \hline 1 & 1 & 1 \\ \hline 2 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_8 = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 0 \\ \hline 1 & 1 & 1 \\ \hline 0 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{368} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 1 \\ \hline 0 & 1 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{76} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline 0 & 0 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 2 \\ \hline 1 & 1 & 1 \\ \hline 2 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 1 & 2 \\ \hline 1 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{684} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 2 \\ \hline 0 & 2 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline 1 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 0 & 2 \\ \hline 0 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{332} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 1 \\ \hline 0 & 1 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 0 \\ \hline 1 & 1 & 1 \\ \hline 0 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

$$s_{424} = \begin{array}{|c|c|c|} \hline 0 & 0 & 0 \\ \hline 0 & 0 & 1 \\ \hline 0 & 2 & 0 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 1 & 1 & 0 \\ \hline 1 & 1 & 1 \\ \hline 2 & 1 & 1 \\ \hline \end{array} \quad \begin{array}{|c|c|c|} \hline 2 & 0 & 2 \\ \hline 1 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array}$$

§2 Tools of Our Research

◇ Polynomials over a Finite Field ◇

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◇ Polynomials over a Finite Field ◇

Let $k (> 1)$ be a power of a prime.

We consider the base set

$$E_k = \{0, 1, \dots, k - 1\}$$

as a finite field (Galois field) $\text{GF}(k)$,

and

express a function defined on E_k as a polynomial over $\text{GF}(k)$.

QUIZ (Easy)

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Given polynomials f, g over $\text{GF}(2)$
(i.e., f, g : Boolean function)

$$f(x, y) = xy + 1$$

$$g(x, y) = xy + x + y$$

QUIZ (Easy)

Given polynomials f, g over $\text{GF}(2)$
(i.e., f, g : Boolean function)

$$f(x, y) = xy + 1$$

$$g(x, y) = xy + x + y$$

QUESTION :

Which is weaker, f or g ?

(with respect to the productive
power of functions)

ANSWER :

g is much weaker than *f*

ANSWER :

g is much weaker than f

Because :

$$f(x, y) = \text{NAND}(x, y)$$

$$g(x, y) = \text{OR}(x, y)$$

QUIZ (Hard)

QUIZ (Hard)

Given polynomials u, v, w over $\mathbf{GF}(3)$

$$u(x, y) = x^2y^2 + xy^2 + x^2y + 2xy + x + y$$

$$v(x, y) = x^2y^2 + xy^2 + x^2y + xy + x + y$$

$$w(x, y) = x^2y^2 + xy^2 + x^2y + 2xy + x + y + 1$$

QUIZ (Hard)

Given polynomials u, v, w over $\text{GF}(3)$

$$u(x, y) = x^2y^2 + xy^2 + x^2y + 2xy + x + y$$

$$v(x, y) = x^2y^2 + xy^2 + x^2y + xy + x + y$$

$$w(x, y) = x^2y^2 + xy^2 + x^2y + 2xy + x + y + 1$$

QUESTION :

**Which is the weakest
among u, v and w ?**

**(with respect to the productive
power of functions)**

ANSWER :

u is the weakest.

ANSWER :

u is the weakest.

Because :

- (1)** $u(x, y)$ generates a minimal clone.
- (3)** $w(x, y)$ is “Webb function” which generates all functions.
- (2)** $v(x, y)$ is inbetween.

Our Hope !!



To characterize polynomials generating minimal clones in terms of the “form of polynomials”

or

To discover “principles” or “rules” that are common to polynomials generating minimal clones



§3 Basic Facts on Minimal Clones

Lemma

A minimal clone is generated by a single function,

i.e., for any minimal clone $C \in \mathcal{L}_k$ there exists $f \in \mathcal{O}_k$ such that $C = \langle f \rangle$

Proposition

For any $f \in \mathcal{O}_k$, suppose f satisfies the following :

- (i) $f \notin \mathcal{J}_k$
- (ii) **For** $\forall h \in \langle f \rangle$, **if** $h \notin \mathcal{J}_k$ **then**
 $f \in \langle h \rangle$

Then f **generates** a minimal clone.

Type Theorem for Minimal Clones

An function f on E_k is *minimal* if

- (i) it generates a minimal clone, and
- (ii) every function from $\langle f \rangle$ whose arity is smaller than the arity of f is a projection

Type Theorem for Minimal Clones

An function f on E_k is *minimal* if

- (i) it generates a minimal clone, and
- (ii) every function from $\langle f \rangle$ whose arity is smaller than the arity of f is a projection

Theorem (I. G. Rosenberg, 1986 ;

“Inspired by Csákány’s work, ... ”)

Every minimal function is classified in one of the following five types:

- (1) Unary function
- (2) Idempotent binary function
- (3) Majority function
- (4) Semiprojection
- (5) If $k = 2^m$, the ternary function $f(x, y, z) := x + y + z$ where $\langle E_k; + \rangle$ is an elementary 2-group

Note :

$f(x_1, \dots, x_n)$: **idempotent**

\iff

$f(x, \dots, x) = x$ **for** $\forall x \in E_k$

♡ Generators of all minimal clones
of type (2) over GF(3) ♡

(Originally from B. Csákány)

$$b_{11} = xy^2$$

$$b_{624} = 2x + 2y$$

$$b_{68} = 2x + 2xy^2$$

$$b_0 = 2x^2y + 2xy^2$$

$$b_{449} = x + y + 2x^2y$$

$$b_{368} = x + y^2 + 2x^2y^2$$

$$b_{692} = x + 2y^2 + x^2y^2$$

$$b_{33} = x + 2x^2y + xy^2$$

$$b_{41} = x^2 + xy^2 + 2x^2y^2$$

$$b_{71} = 2x^2 + xy^2 + x^2y^2$$

$$b_{26} = 2x + x^2 + 2xy + 2x^2y$$

$$b_{37} = 2x + 2x^2 + xy + 2x^2y$$

$$b_{17} = 2x + x^2 + 2xy^2 + 2x^2y^2$$

$$b_{38} = 2x + 2x^2 + 2xy^2 + x^2y^2$$

$$b_{10} = xy + 2x^2y + 2xy^2 + 2x^2y^2$$

$$\begin{aligned}
b_{20} &= 2xy + 2x^2y + 2xy^2 + x^2y^2 \\
b_{43} &= x + xy + 2x^2y + xy^2 + 2x^2y^2 \\
b_{53} &= x + 2xy + 2x^2y + xy^2 + x^2y^2 \\
b_{35} &= x + xy + x^2y + 2xy^2 + 2x^2y^2 \\
b_{42} &= x + 2xy + x^2y + 2xy^2 + x^2y^2 \\
b_{530} &= x + y + y^2 + 2x^2y + 2x^2y^2 \\
b_{125} &= x + y + 2y^2 + 2x^2y + x^2y^2 \\
b_{116} &= x + y + xy + 2y^2 + 2xy^2 \\
b_{528} &= x + y + 2xy + y^2 + 2xy^2 \\
b_{206} &= x + 2y + y^2 + x^2y + 2x^2y^2 \\
b_{287} &= x + 2y + 2y^2 + x^2y + x^2y^2 \\
b_{215} &= x + 2y + 2xy + y^2 + xy^2 \\
b_{286} &= x + 2y + xy + 2y^2 + xy^2 \\
b_{122} &= y + x^2 + 2y^2 + 2x^2y + xy^2 \\
b_{557} &= y + 2x^2 + y^2 + 2x^2y + xy^2 \\
b_{16} &= 2x + x^2 + xy + 2x^2y + x^2y^2 \\
b_{47} &= 2x + 2x^2 + 2xy + 2x^2y + 2x^2y^2 \\
b_{178} &= 2x + 2y + x^2 + xy + y^2 \\
b_{290} &= 2x + 2y + 2x^2 + 2xy + 2y^2 \\
b_{40} &= x^2 + xy + 2x^2y + 2xy^2 + x^2y^2
\end{aligned}$$

$$\begin{aligned}
b_{80} &= 2x^2 + 2xy + 2x^2y + 2xy^2 + 2x^2y^2 \\
b_{364} &= x^2 + xy + y^2 + 2x^2y + 2xy^2 \\
b_{728} &= 2x^2 + 2xy + 2y^2 + 2x^2y + 2xy^2 \\
b_{448} &= x + y + xy + x^2y + xy^2 + 2x^2y^2 \\
b_{458} &= x + y + 2xy + x^2y + xy^2 + x^2y^2 \\
b_{205} &= x + 2y + xy + y^2 + xy^2 + x^2y^2 \\
b_{296} &= x + 2y + 2xy + 2y^2 + xy^2 + 2x^2y^2 \\
b_{188} &= 2x + 2y + x^2 + 2xy + y^2 + 2x^2y^2 \\
b_{280} &= 2x + 2y + 2x^2 + xy + 2y^2 + x^2y^2 \\
b_8 &= 2x + x^2 + xy + x^2y + xy^2 + x^2y^2 \\
b_{36} &= 2x + 2x^2 + 2xy + x^2y + xy^2 + 2x^2y^2 \\
b_{179} &= 2x + 2y + x^2 + y^2 + x^2y + 2xy^2 + x^2y^2 \\
b_{281} &= 2x + 2y + 2x^2 + 2y^2 + x^2y + 2xy^2 + 2x^2y^2
\end{aligned}$$

§4 More Properties of Minimal Clones

For $f, g \in \mathcal{O}_k$ we write

$$f \rightarrow g \quad \text{if} \quad g \in \langle f \rangle$$

(Note : Binary relation \rightarrow on \mathcal{O}_k is a quasi-order)

Lemma

Let $f \in \mathcal{O}_k^{(2)}$ be an essentially binary function. If f satisfies (1) and (2) then f is minimal.

(1) f is idempotent

(2) For any $g \in \mathcal{O}_k^{(m)}$ ($m \geq 2$)

if $f \rightarrow g$ then g is a projection

or

$$g \rightarrow f$$

Let $m \geq 3$ and $g \in \mathcal{O}_k^{(m)}$. We say

g is a *quasi-projection*

if g becomes a projection whenever two arguments of g are identified, i.e., if $g(x_1, \dots, x_i, \dots, x_i, \dots, x_m)$ is always a projection.

Lemma

Let $f \in \mathcal{O}_k^{(2)}$ be a binary idempotent function. Then

f is minimal \iff (1) and (2) hold

(1) For $\forall g \in \mathcal{O}_k^{(2)} \setminus \mathcal{J}_k$

if $f \rightarrow g$ then $g \rightarrow f$

(2) For $\forall m \geq 3$ and $\forall g \in \mathcal{O}_k^{(m)} \setminus \mathcal{J}_k$

if $f \rightarrow g$ then g is not a quasi-projection.

For $f \in \mathcal{O}_k^{(2)}$ let $\Gamma_f^{(x,y)}$ be :

$$\{ f(f(x,y), x), f(f(x,y), y), f(x, f(x,y)), \\ f(y, f(x,y)), f(f(x,y), f(y,x)) \}$$

Then let $\Gamma_f = \Gamma_f^{(x,y)} \cup \Gamma_f^{(y,x)}$

Lemma

Let $f \in \mathcal{O}_k^{(2)}$ be a binary idempotent function which is not a projection. Suppose that, for any $\gamma \in \Gamma_f$, one of the following holds:

$$\gamma(x,y) \approx f(x,y) \quad \text{or} \quad \gamma(x,y) \approx f(y,x)$$

Then f is minimal

(Here, by $h_1(x,y) \approx h_2(x,y)$ we mean $h_1(x,y) = h_2(x,y)$ for all $(x,y) \in E_k^2$)

§5 Polynomials generating Minimal Clones

Here we consider only

“Binary Idempotent Minimal Functions”

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“Binary Idempotent Minimal Functions”

Our Strategy :

Step 1 : Take arbitrary $f(x, y) \in \mathcal{O}_3^{(2)}$
from Csákány’s list.

Step 2 : Search for a polynomial
 $g(x, y) \in \mathcal{O}_k^{(2)}$ for $k \geq 3$ whose
counterpart for $k = 3$ is $f(x, y)$

Step 3 : Examine if g is minimal.

(1) Linear Polynomials

For $k = 3$, the binary linear polynomial

$$f(x, y) = 2x + 2y$$

is minimal.

For $k = 5$, the binary linear polynomial

$$f(x, y) = 2x + 4y$$

is minimal, whose Cayley table is :

$$\underline{f(x, y) = 2x + 4y}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|----------|----------|----------|----------|----------|
| 0 | 0 | 4 | 3 | 2 | 1 |
| 1 | 2 | 1 | 0 | 4 | 3 |
| 2 | 4 | 3 | 2 | 1 | 0 |
| 3 | 1 | 0 | 4 | 3 | 2 |
| 4 | 3 | 2 | 1 | 0 | 4 |

Theorem (Á. Szendrei)

Let k be a prime. Let $f(x, y)$ be a binary linear polynomial on E_k

Then f is *minimal* if and only if

$$f(x, y) = ax + (k + 1 - a)y$$

for some $1 < a < k$

Moreover, all such linear polynomials generate the same minimal clone.

(2) Monomials

Consider $x^s y^t$ for $1 \leq s \leq t < k$

For $k = 3$, $f(x, y) = x y^2$ is minimal.

For $k = 5$,

$$f(x, y) = x y^4$$

is minimal whose Cayley table is :

$$\underline{f(x, y) = x y^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|----------|----------|----------|----------|----------|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 | 1 |
| 2 | 0 | 2 | 2 | 2 | 2 |
| 3 | 0 | 3 | 3 | 3 | 3 |
| 4 | 0 | 4 | 4 | 4 | 4 |

Theorem

Let k be a prime.

(1) $x y^{k-1}$ is a minimal function.

(2) For any $1 < s < k - 1$,
 $x^s y^{k-s}$ is not a minimal
function.

(3) More Examples

(3) More Examples

Recall !!

Our Strategy :

Step 1 : Take arbitrary $f(x, y) \in \mathcal{O}_3^{(2)}$
from Csákány's list.

Step 2 : Search for a polynomial
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counterpart for $k = 3$ is $f(x, y)$

Step 3 : Examine if g is minimal.

(3) More Examples

(3-1) Example A

Step 1 :

$$k = 3 : x + y + 2xy^2$$

(3) More Examples

(3-1) Example A

Step 1 :

$$k = 3 : x + y + 2xy^2$$

Step 2 :

$$k \geq 3 : \quad ? ? ?$$

(3) More Examples

(3-1) Example A

Step 1 :

$$k = 3 : x + y + 2 x y^2$$

Step 2 :

$$k \geq 3 : x + y + (k - 1) x y^{k-1}$$

(3) More Examples

(3-1) Example A

Step 1 :

$$k = 3 : x + y + 2xy^2$$

Step 2 :

$$k \geq 3 : x + y + (k - 1)xy^{k-1}$$

Example. $k = 5$:

$$\underline{f(x, y) = x + y + 4xy^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|---|---|---|---|---|
| 0 | 0 | 1 | 2 | 3 | 4 |
| 1 | 1 | 1 | 2 | 3 | 4 |
| 2 | 2 | 1 | 2 | 3 | 4 |
| 3 | 3 | 1 | 2 | 3 | 4 |
| 4 | 4 | 1 | 2 | 3 | 4 |

(3-2) Example B

Step 1 :

$$k = 3 : x + 2y^2 + x^2 y^2$$

(3-2) Example B

Step 1 :

$$k = 3 : x + 2y^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 : \quad ? \quad ? \quad ?$$

(3-2) Example B

Step 1 :

$$k = 3 : x + 2y^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 :$$

$$x + (k - 1)y^{k-1} + x^{k-1} y^{k-1}$$

(3-2) Example B

Step 1 :

$$k = 3 : x y^2 + 2 x^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 :$$

$$x + (k - 1)y^{k-1} + x^{k-1} y^{k-1}$$

Example. $k = 5 :$

$$\underline{f(x, y) = x + 4y^4 + x^4 y^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|---|---|---|---|---|
| 0 | 0 | 4 | 4 | 4 | 4 |
| 1 | 1 | 1 | 1 | 1 | 1 |
| 2 | 2 | 2 | 2 | 2 | 2 |
| 3 | 3 | 3 | 3 | 3 | 3 |
| 4 | 4 | 4 | 4 | 4 | 4 |

(3-3) Example C

Step 1 :

$$k = 3 : x y^2 + 2 x^2 + x^2 y^2$$

(3-3) Example C

Step 1 :

$$k = 3 : x y^2 + 2 x^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 : \quad ? ? ?$$

(3-3) Example C

Step 1 :

$$k = 3 : x y^2 + 2 x^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 :$$

$$x y^{k-1} + (k-1)x^{k-1} + x^{k-1} y^{k-1}$$

(3-3) Example C

Step 1 :

$$k = 3 : x y^2 + 2 x^2 + x^2 y^2$$

Step 2 :

$$k \geq 3 :$$

$$x y^{k-1} + (k-1)x^{k-1} + x^{k-1} y^{k-1}$$

Example. $k = 5 :$

$$\underline{f(x, y) = x y^4 + 4 x^4 + x^4 y^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 4 | 1 | 1 | 1 | 1 |
| 2 | 4 | 2 | 2 | 2 | 2 |
| 3 | 4 | 3 | 3 | 3 | 3 |
| 4 | 4 | 4 | 4 | 4 | 4 |

(3-4) Example D

A bit more difficult example !

(3-4) Example D

Step 1 :

$$k = 3 : 2x + 2xy^2$$

(3-4) Example D

Step 1 :

$$k = 3 : 2x + 2xy^2$$

Step 2 :

$$k \geq 3 : \quad ? ? ?$$

(3-4) Example D

Step 1 :

$$k = 3 : 2x + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1)x + (k - 1)xy^{k-1}$$

is not good !

(3-4) Example D

Step 1 :

$$k = 3 : 2x + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1)x + 2xy^{k-1}$$

(3-4) Example D

Step 1 :

$$k = 3 : 2x + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1)x + 2xy^{k-1}$$

Example. $k = 5 :$

$$\underline{f(x, y) = 4x + 2xy^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 4 | 1 | 1 | 1 | 1 |
| 2 | 3 | 2 | 2 | 2 | 2 |
| 3 | 2 | 3 | 3 | 3 | 3 |
| 4 | 1 | 4 | 4 | 4 | 4 |

(3-5) Example E

This example requires better skill
even to find
a candidate of generalization !!

(3-5) Example E

Step 1 :

$$k = 3 : 2x^2y + 2xy^2$$

(3-5) Example E

Step 1 :

$$k = 3 : 2x^2y + 2xy^2$$

Step 2 :

$$k \geq 3 : \quad ? ? ?$$

(3-5) Example E

Step 1 :

$$k = 3 : 2x^2y + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1)x^{k-1}y + (k - 1)xy^{k-1}$$

is not good !

(3-5) Example E

Step 1 :

$$k = 3 : 2 x^2 y + 2 x y^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1) x^{k-1} y + (k - 1) x y^{k-1}$$

is not good !

$$2 x^{k-1} y + (k - 1) x y^{k-1}$$

is also not good !!

(3-5) Example E

Step 1 :

$$k = 3 : 2x^2y + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1) \sum_{i=1}^{k-1} x^{k-i} y^i$$

(3-5) Example E

Step 1 :

$$k = 3 : 2x^2y + 2xy^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1) \sum_{i=1}^{k-1} x^{k-i} y^i$$

Step 3 (Sketch of Proof):

◇ $f(x, x) = x$ since $(k - 1)^2 = 1$

◇ For $x \neq y$, let $D = \sum_{i=1}^{k-1} x^{k-i} y^i$

Then

$$xy^{-1}D = D$$

Hence $xD = yD$ which implies $D = 0$

Therefore $f(x, y) = x$ if $x = y$

and $f(x, y) = 0$ if $x \neq y$

It is easy to see that f is minimal

(3-5) Example E

Step 1 :

$$k = 3 : 2 x^2 y + 2 x y^2$$

Step 2 :

$$k \geq 3 :$$

$$(k - 1) \sum_{i=1}^{k-1} x^{k-i} y^i$$

Example. $k = 5 :$

$$\underline{f(x, y) = 4 x^4 y + 4 x^3 y^2 + 4 x^2 y^3 + 4 x y^4}$$

| $x \backslash y$ | 0 | 1 | 2 | 3 | 4 |
|------------------|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 |
| 2 | 0 | 0 | 2 | 0 | 0 |
| 3 | 0 | 0 | 0 | 3 | 0 |
| 4 | 0 | 0 | 0 | 0 | 4 |

§ ∞ More and More Examples

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You will hear at the conference
celebrating
the 80th Birthday of
Béla Csákány !!

§ ∞ More and More Examples

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Thank you